A Study of Soft Matrices

loia: la.

Ja"

Department of Mathematics and Statistics Faculty of Basic and Applied Sciences International Islamic University, Islamabad, Pakistan. 20L6

Accession No $\frac{7H-16063}{4}$

 M^* MS 512.896 KHS

1. Matrices
2. Mathematical analysis

A Study of Soft Matrices

 \mathbf{r}

 $\frac{1}{2\sqrt{3}}\frac{1}{2(1-\sqrt{3})}\frac{1}{\sqrt{3}}$

 $\ddot{\bullet}$

こんじょう

ţ. $\pmb{\ddagger}$

こうしょう

By **Azhar Rauf Khan**

Supervised by Dr. Tahir Mahmood

Department of Mathematics and Statistics Faculty of Basic and Applied Sciences International Islamic University, Islamabad, Pakistan. 2016

A Study of Soft Matrices

By **Azhar Rauf Khan**

寥

 $\mathbf{1}$ \mathbf{r}

籀

A Dissertation **Submitted in the Partial Fulfillment of the Requirements for the Degree of MASTER OF SCIENCE In MATHEMATICS**

> Supervised by Dr. Tahir Mahmood

Pepartment of Mathematics and Statistics Faculty of Basic and Applied Sciences International Islamic University, Islamabad, Pakistan. 2016

Certificate

A Study of Soft Matrices

By

Azhar Rauf Khan

A DISSERTATION SUBMITTED IN THE PARTIAL FULFILLMENT OF THE REQUIREMENTS FOR THE DEGREE OF THE MASTER OF SCIENCE IN MATHEMATICS

We accept this dissertation as conforming to the required standard.

 $2.$

 $\mathbf{1}$.

 $\overline{3}$.

Prof. Dr. Muhammad Shabir (External Examiner)

> 'ahir Mahmood Supervisor)

 $22/06$ Prof. Dr. M. Arshad Zia (Internal Examiner)

 $\overline{\mathcal{L}}$

 $89/16$ Prof. Dr. M. Arshad Zia (Chairman)

Department of Mathematics and Statistics Faculty of Basic and Applied Sciences International Islamic University, Islamabad, Pakistan. 2016

Forwarding sheet by Research Supervisor

This thesis entitled "A Study of Soft Matrices" submitted by Azhar Rauf Khan [Reg. No. 17S-FBAS/MSMA /S-14) is partial fulfillment of MS degree in Mathematics is completed under my guidance and supervision. I am satisfied with the quality of his research work and allow him to submit this thesis for further process to graduate with Master of, Science degree from Department of Mathematics and Statistics, Faculty of Basic and Applied Sciences as per International Islamic University, Islamabad rules and regulations.

Date:

 ϵ

i $\hat{\mathbf{z}}_{\hat{\eta}}$

ft. F

Dr. Tahir Mahmood

l}

I

Assistant professor Department of Mathematics and Statistics Faculty of Basic and Applied Sciences International Islamic University, Islamabad, Pakistan.

DECLARATION

I hereby, declare, that this thesis neither as a.whole nor as a part of it has been copied out from any source. It is further declared that I have prepared this thesis $\sum_{n=1}^{\infty}$ entirely on the basis of my personal efforts made under the sincere guidance of my kind supervisor Dr. Tahir Mahmood. No portion of the work, presented in this thesis, has been submitted in the support of any application for any degree or qualification of this or any other institute of learning.

Azhar RaufKhan

,A ry

 \mathcal{F}

MS Mathematics Reg. No. 175-FBAS/MSMA/S-14. Department of Mathematics and Statistics Faculty of Basic and Applied Sciences International Islamic University, Islamabad, Pakistan.

Dedicated To **My Loving Parents, Respectful Teachers And My Friends.**

 \mathbb{P}^{\bullet}

Acknowledgements

First and foremost, I owe profound gratitude to Almighty ALLAH, the most merciful and most compassionate, the creator of this Universe for Man to explore it, and who bestowed upon-me strength and ability to complete my research work. I offer my humblest and sincerest words of thahks to HOLY PROPHET MUHAMMAD (PEACE BE UPON HIM) who is, forever, a torch of guidance and knowledge for all the human beings.

. I would like to express my sincere gratitude to Dr. Tahir Mahmood, my teacher and supervisor, for continuous support and guidance of my MS. study and research. He always encouraged me and criticized positively, gave me his time whenever and wherever I needed. This research work would not have been possible without their kind support and the creative abilities. In spite of his extremely busy schedule, he always used to iake their precious time for me. In short, he proved themselves to be a perfect model of professionalism, understanding and commitment to the subject and to his students. I could not have imagined having a better advisors and mentor for my MS. study.

Dr. Muhammad Irfan Ali my teacher his guidance and support has given me much strength and courage to complete my research. He has given a right direction to my work and his instructive idea proves very fruitful for me in the completion of my research.

I am highly grateful to Dr. Muhammad Arshad Zia, Chairman, Department of Mathematics and statistics, International Islamic University Islamabad, for the provision of all possible facilities and for their full cooperation.

It would have not been possible for me to reach at this stage without guidance of all my teachers, especially Dr. Muhammad Arshad Zia, Dr. Rahmat Ellahi, Dr. Nasir Ali, Dr. Ahmad Zeeshan, Dr. Ahmer Mahmood, Dr. Tariq Javed, Dr. Nayyar Mehmood and Mr. Niaz Ahmad. I owe my gratitude to all of them.

I am very thankful to all members of my family, for their support, affection, encouragement attribute and my success to my mother whose love, prayers and guidance always gave me a ray of hope in the darkness of desperation. I am very much indebted to my brothers and sister for their love and support.

In my entire journey, I am very grateful to Dr. Saqib Hussain, Rasib Awan, Irfan Nazir, Mohsin ali khan, Asim Zamir, Muhammad Mubashar Abbasi, Abdullah, Abid Rehman, Qaisir khan, Faisal Mehmood, Azmat Hussain Turi and Muhammad Siddique, for their support, help and encouragement. I owe thanks to all of them.

Azhar RaufKhan

a

g' \downarrow

'+rr কু

Contents

 $\hat{\mathbf{E}}$

愿

E.

Preface

 $\bar{\mathbf{e}}$

 $\tilde{\mathcal{F}}$

j:(v

In this thesis, the center of discussion is some Iimitations of soft set matrices and its uses. The soft sets concept was expressed by Molodtsov in 1999 [17]. This concept is used to solve some complications in the fields of economics, engineering and environment because all these areas have some distinctive uncertainties regarding these problems. The concept of soft set is applied in fuzzy sets, intuitionistic fuzzy sets, vague set, interval mathematics and rough sets. In this thesis, some discussion is also done on matrices, which have ^a significant role in the vast field of engineering, economics and science. But, the old theory related to matrices is failed in solving the uncertainties, which are caused due to inaccurate circumstances. Matrices have different properties which include: commutative law, associative law and distributive law.

ffi

In the study, the idea of soft sets is described by linking an advantageous method with soft matrices. This study also involves the Naim Cagmanis and S. Enginoslu [5] research which highlights the usage of soft set theory in more precise manner. He describes the different dimensions of its applications. Initially, with the help of rough sets, he gave the theory of soft sets in decision making problems. Xiao et. all 127) had done a research highlighting business competitive capacity based on soft sets. Maji et al, [13] defined the fuzzy set, as the time passes a lot of work has done in fuzzy soft set. The definition of soft group was given by Aktas and Cagman [1]. They also made a comparison between soft sets to the rough soft sets and fuzzy soft sets, Subsequently, many other researchers have done a lot of work on this concept and gave many other theories related to the soft sets. Roy and maji [25] have also done some work on the applications and decision making problem. Majumdar [16] introduced the reduction of fuzzy soft set and then examine a decision making problem by fuzzy soft sets. The theory of the Rough sets is explained by Pawlak [23] for the analysis of the data possibly with inconsistent information. This theory has been used in many fields such as beauty contest, pattern recognition confliict analysis and switching circuits.

In the light of above mentioned facts, we indicates some limitations of the products of soft matrices given by Naim Cagman [5]. We pointed out that the products of soft matrices are not binary. It does not satisfy many laws which include Closure law, associative law and distributive law. Keeping in view this drawback in this thesis we have introduced new products of soft matrices, which are binary. We have also shown that accociative laws and distributive laws also holds.

 \mathbf{i}

Structure of the Thesis

The thesis is organized chapter wise as follows:

Chapter 1:

This chapter is introductory and sets up the background for the problems taken in the thesis. Semirings, Soft Sets, Soft-Union-lntersection Sum, Soft-Union-Intersection Product and related results are discussed.

Chapter 2:

In this chapter the article "Soft matrix theory and its decision making" is reviewed.

Chapter 3:

 $\overline{\mathbf{C}}$

 \mathcal{L}

 \sim :s"

In this chapter, keeping in view the drawbacks and limitation such as the products of soft matrices defined in the paper reviewed are not binary and that associative and distributive laws are not satisfied, we improved the products of soft matrices and named them B-products of soft matrices. It is also shown that the defined products are binary. Further it is also shown that these products now Satisfy the associative laws and distributive laws as well.

 $\ddot{}$ I

И

 \mathbf{I}

Ĩ

4

.t

Chapter ¹

'ui

優

Preliminaries

This chapter provides the essential definitions and preliminary results, which are useful for our subsequent chapters. For undefined terms and notions we refer to $([1], [2], [3],$ $[4], [5], [8], [10], [14], [16], [15], [17], [23], [25], [27].$

1.1 Semigroups

Let S be a non-empty set and "*" be a binary operation on S. Then $(S, *)$ is called a semigroup if this operation is associative, that is

 $a*(b*c)=(a*b)*c$ for all $a,b,c\in S$.

A semigroup $(S, *)$ is called *commutative* if.

 $a * b = b * a$ for all $a, b \in S$.

1.1.1 Definition

Let $(S, *)$ be a semigroup. If there exists an element $e \in S$ such that

 $a*e=e*a=a$ for all $a\in S$,

then e is called the *identity element* in S and $(S, *)$ is called a *monoid*.

An element $x \in S$ is called idempotent if $x * x = x$. If every element of S is idempotent then we say that S is idempotent.

Usually instead of writing $(S, *)$ we write S & instead of writing $x * y$ we write xy , for all $x,y \in S$.

 $\mathbf 1$

1.1.2 Examples

 $\mathcal{D}% _{M_{1},M_{2}}^{\alpha,\beta}(\varepsilon)$

R

Ñ

- 1. $(N, +)$ is a semigroup.
- 2. Let $S = \{a_1, a_2, a_3, \dots\}$ such that $*$ be defined on S by $a_i * a_j = a_i$. Then $(S,*)$ is a semigroup.
- 3. $(N_0, +)$ is a Monoid, where $N_0 = N \cup \{0\}$
- 4. (\mathbb{Z}, \cdot) is a Monoid.
- 5. $\{0, 1\}$ is a monoid under ".".
- 6. For any set X ; $(P(X), \cup)$ and $(P(X), \cap)$ are monoids.

L.2 Semirings

A semiring is an algebraic system consisting of a non-empty set R together with two binary operations called "addition" and "multiplication" (denoted by "+" and ".", respectively) such that $(R,+)$ and $(R,.)$ are semigroups and multiplication distributes over addition from both sides, that is

 $a \cdot (b+c) = a \cdot b + a \cdot c$, and $(b+c) \cdot a = b \cdot a + c \cdot a$

for all $a, b, c \in R$.

1.3 Soft Sets

Soft set theory was introduced by D. Molodtsov [17]. It is a new approach for the real world problems in the field of economics, engineering, management etc. Molodtsov's soft set theory was proposed for dealing with ambiguity. He also defined some operations for soft set theory.

1.3.1 Definition [17]

Let U be an initial universe ; E be the set of all possible parameters under consideration with respect to U and A be a subset of E. Then a pair (F, A) is called a soft set over U, where F is a mapping given by, $F : A \to \mathcal{P}(U)$.

For $e \in A$, $F(e)$ may be considered as the set of e-approximate elements of the soft set (F, A) .

-Parameters are often attributes, characteristics, or properties of objects in soft sets. For example big, airy, tall, cool, hot, wooden, expensive, cheap etc.

.F\

G

In other words, a soft set over U is a parameterized family of subsets of the universe U. For $e \in A$, $F(e)$ may be considered as the set of e-approximate elements of the soft set (F, A) .

1.3.2 Definition [15]

For two soft sets (F, A) and (G, B) over a common universe U, we say that (F, A) is a soft subset of (G, B) if

1. $A \subseteq B$ and

2. $F(e) \subseteq G(e)$ for all $e \in A$.

We write $(F, A)\tilde{\subset} (G, B)$.

In this case (G, B) is said to be a soft super set of (F, A) .

1.3.3 Definition [15]

Two soft sets (F, A) and (G, B) over a common universe U are said to be soft equal if (F, A) is a soft subset of (G, B) and (G, B) is a soft subset of (F, A) .

1.3.4 Definition $\overline{2}$

Let U be an initial universe set, E be the set of parameters, and $A \subseteq E$.

- 1. (F, A) is called a relative null soft set (with respect to the parameter set A), denoted by \emptyset_A , if $F(a) = \emptyset$ for all $a \in A$.
- 2. (G, A) is called a relative whole soft set (with respect to the parameter set A), denoted by \mathfrak{U}_A , if $G(a) = U$ for all $a \in A$..

The relative whole soft set with respect to the set of parameters E is called the absolute soft set over U and denoted by \mathfrak{U}_E . In a similar way, the relative null soft set with respect to E is called the *null soft set* over U and is denoted by \varnothing _E.

We shall denote by \mathcal{O}_\emptyset the unique soft set over U with an empty parameter set, which is called the *empty soft set* over U. Note that \mathcal{O}_{q} and \mathcal{O}_{A} are different soft sets over U and $\mathcal{O}_{\emptyset} \widetilde{\subset} \mathcal{O}_{A} \widetilde{\subset} (F, A) \widetilde{\subset} \mathfrak{U}_{A} \widetilde{\subset} \mathfrak{U}_{E}$ for all soft set (F, A) over U.

1.3.5 Definition [2]

Extended union of two soft sets (F, A) and (G, B) over the common universe U is the soft set (H, C) , where $C = A \cup B$ and for all $e \in C$,

 \mathbf{G}

 $\tilde{\bm{v}}$

$$
H(e) = \begin{cases} F(e) & \text{if } e \in A - B \\ G(e) & \text{if } e \in B - A \\ F(e) \cup G(e) & \text{if } e \in A \cap B \end{cases}
$$
\n
$$
\text{We write } (F, A) \cup_{\mathcal{E}} (G, B) = (H, C).
$$

1.3.6 Definition [2]

Let (F, A) and (G, B) be two soft sets over the same universe U, such that $A \cap B \neq \emptyset$. The restricted union of (F, A) and (G, B) is denoted by $(F, A) \cup_{\mathcal{R}} (G, B)$ and is defined as $(F, A) \cup_{\mathcal{R}} (G, B) = (H, C)$, where $C = A \cap B$ and for all $e \in C$, $H(e) = F(e) \cup G(e)$. If $A \cap B = \emptyset$, then $(F, A) \cup_{\mathcal{R}} (G, B) = \emptyset_{\emptyset}$.

1.3.7 Definition $[2]$

The extended intersection of two soft sets (F, A) and (G, B) over a common universe U, is the soft set (H, C) where $C = A \cup B$ and for all $e \in C$,

 $\int F(e)$ if $e \in A-B$ $H(e)=\left\{ \begin{array}{ll} G\left(e\right)&\quad if\;e\in B-A \end{array} \right.$ $\left(\begin{array}{cc} F\left(e\right) \cap G\left(e\right) & if\; e\in A\cap B \end{array} \right)$ We write $(F, A) \cap_{\mathcal{E}} (G, B) = (H, C)$

1.3.8 Definition [2]

Let (F, A) and (G, B) be two soft sets over the same universe U such that $A \cap B \neq \emptyset$. The restricted intersection of (F, A) and (G, B) is denoted by $(F, A) \cap_{\mathcal{R}} (G, B)$ and is defined as $(F, A) \cap_{\mathcal{R}} (G, B) = (H, A \cap B)$, where $H(e) = F(e) \cap G(e)$ for all $e \in A \cap B$. If $A \cap B = \emptyset$ then $(F, A) \cap_{\mathcal{R}} (G, B) = \emptyset_{\emptyset}$.

1.3.9 Definition [2]

Let (F, A) and (G, B) be two soft sets over the same universe U such that $A \cap B \neq \emptyset$. The restricted difference of (F, A) and (G, B) is denoted by $(F, A) \sim_{\mathcal{R}} (G, B)$ and is defined as $(F, A) \smile_{\mathcal{R}} (G, B) = (H, A \cap B)$, where $H(e) = F(e) - G(e)$ for all $e \in A \cap B$.

If $A \cap B = \emptyset$ then $(F, A) \cup_R (G, B) = \emptyset_{\emptyset}$.

1.3.L0 Definition [2]

The complement of-a soft set (F, A) is denoted by $(F, A)^c$ and is defined by $(F, A)^c$ $(A)^c = (F^c, A)$ where $F^c : A \to P(U)$ is mapping given by $F^c(e) = U - F(e)$ for all $e \in A$

il ir to I ${^{\circ}}$,t { I

 $\frac{1}{2}$

1. Preliminaries

 \hat{v}

s.

ie clearly $(F, A)^c = \mathfrak{U}_A \smile_{\mathcal{R}} (F, A)$ and $((F, A)^c)^c = (F, A)$.

1.3.11 Definition [15]

Let (F, A) and (G, B) be any two soft sets over a common universe U. Then the basic intersection of two soft sets (F, A) and (G, B) is defined as the soft set (H, C) = $(F, A) \wedge (G, B)$, where $C = A \times B$, and $H(a, b) = F(a) \cap G(b)$ for all $(a, b) \in A \times B$.

L.3.L2 Definition [15]

Let (F, A) and (G, B) be any two soft sets over a common universe U. Then the basic union of two soft sets (F, A) and (G, B) is defined as the soft set (H, C) = $(F, A) \vee (G, B)$, where $C = A \times B$, and $H(a, b) = F(a) \cup G(b)$ for all $(a, b) \in A \times B$.

1.3.13 Theorem

Let (F, A) and (G, B) be two soft sets over the same universe \vec{U} such that $A \cap B \neq \emptyset$. Then

- (1) $((F, A)\cup_R (G, B))^c = (F, A)^c \cap_R (G, B)^c$
- (2) $((F, A)\cap_R (G, B))^c = (F, A)^c \cup_R (G, B)^c$

1.3.14 Distributive Laws for Soft Sets

In this section, we discuss distributive laws on the collection of soft sets. It is interesting to see that the equality does not hold in each and every case. We see the improperness in some assertions and counter example is given to show it. Let U be an initial universe and E be the set of parameters then we denote the collections of soft set as follows.

 $\mathcal{SS}(U)^E$: The collection of all soft sets defined over U.

 $\mathcal{SS}(U)_A$: The collection of all those soft sets defined over U with a fixed parameters set A.

1.3.L5 Proposition [3]

Let (F, A) be a soft set over the universe set U.

(1) $(F, A)\alpha(F, A) = (F, A)$ for all $\alpha \in {\cap_{\mathcal{R}}}, \cup_{\mathcal{R}}$

(2) $(F, A) \cap_{\mathcal{R}} \emptyset_A = \emptyset_A$

(3) $(F, A) \cup_{\mathcal{R}} \emptyset_A = (F, A)$

$$
(4) (F, A) \cap_{\mathcal{R}} \mathfrak{U}_A = (F, A)
$$

(5) $(F, \tilde{A}) \cup_R \mathfrak{U}_A = \mathfrak{U}_A$

Proof. Straightforward \blacksquare

 \widehat{E}'

 \mathbf{F}

1.3.16 Remark [3]

Let $\alpha, \beta \in \{ \cup, \cap, \cup_{\varepsilon}, \cap_{\varepsilon} \}.$ Then

$$
(F, A) \alpha ((G, B) \beta (H, C)) = ((F, A) \alpha (G, B)) \beta ((F, A) \alpha (H, C))
$$

 \Box holds when we have 1 otherwise 0 in Table 2.

Table 2 shows that, if $\alpha, \beta \in \{\cup_{\mathcal{R}}, \cap_{\mathcal{R}}, \cup_{\varepsilon}, \cap_{\varepsilon}\},\$ then there are sixteen combinations in all, there are four combinations in which $\alpha = \beta$ and for eight combination equality $(F, A) \alpha((G, B) \beta(H, C)) = ((F, A) \alpha(G, B)) \beta((F, A) \alpha(H, C))$ will holds. Proofs in the case where equality holds can^{the} followed by definitions of respective operations. For four remaining α and β this equality does not hold. To show this we have following example

L.3.LT Example [3]

Let U be the set of sample designs and E be the set of available colors for dresses in a boutique,

 $U = \{S_1, S_2, S_3, S_4, S_5, S_6, S_7, S_8\}$ $E = {Red, Green, Blue, Yellow, Black, White, Pink}.$ Suppose that $A = {Red, Green, Blue, White}, B = {Green, Blue, Yellow, Black}$ and $C = \{\text{Blue, Yellow, White, Pink}\}.$ Let $(F, A), (G, B)$ and (H, C) be the soft sets over U, which are defined as follows: $F(\text{Red}) = \{S_1, S_2, S_3, S_4\};$ $F(\text{Blue}) = \{S_1, S_2, S_4, S_7\};$ $G(\text{Green}) = \{S_4, S_5, S_6, S_8\};$ $F(\text{Green}) = \{S_3, S_4, S_5, S_6\};$ $F(\text{White}) = \{S_2, S_3, S_4\}.$ $G(\text{Blue}) = \{S_1, S_2, S_3, S_4\};$ $G(Yelllow) = \{S_4, S_5, S_6, S_7, S_8\};$ $G(Black) = \{S_1, S_2, S_4, S_7\}$ and $H(\text{Blue}) = \{S_3, S_4, S_7, S_8\};$ $H(\text{White}) = \{S_2, S_4, S_6, S_8\};$ Let $H(Yellow) = \{S_4, S_5, S_7\};$ $H(\text{Pink}) = \{S_2, S_3, S_5, S_7\}.$

ł

1. Preliminaries

e)
C

.
براني.
التجار

 $(F, A) \cup_{\varepsilon} ((G, B) \cup_{\mathcal{R}} (H, C)) = (I, A \cup (B \cap C));$ $((F, A) \cup_{\varepsilon} (G, B)) \cup_{\mathcal{R}} ((F, A) \cup_{\varepsilon} (H, C)) = (J, (A \cup B) \cap (A \cup C));$ $(F, A) \cup_{\varepsilon} ((G, B) \cup_{\mathcal{R}} (H, C)) = (K, A \cup (B \cap C))$ $((F, A) \cup_{\varepsilon} (G, B)) \cup_{\mathcal{R}} ((F, A) \cup_{\varepsilon} (H, C)) = (L, (A \cup B) \cap (A \cup C));$ $(F, A) \cup_{\varepsilon} ((G, B) \cup_{\varepsilon} (H, C)) = (M, A \cup (B \cup C));$ $((F, A) \cup_{\varepsilon} (G, B)) \cup_{\varepsilon} ((F, A) \cup_{\varepsilon} (H, C)) = (\dot{N}, (A \cup B) \cup (B \cup C));$ $(F, A) \cup_{\varepsilon} ((G, B) \cup_{\varepsilon} (H, C)) = (O, A \cup (B \cup C));$ $((F, A) \cup_{\varepsilon} (G, B)) \cup_{\varepsilon} ((F, A) \cup_{\varepsilon} (H, C)) = (P, (A \cup B) \cup (B \cup C)).$ Then $I(\text{Red}) = \{S_1, S_2, S_3, S_4\};$ $I(\text{Blue}) = \{S_1, S_2, S_3, S_4, S_7, S_8\};$ $I(\text{White}) = \{S_2, S_3, S_4\}.$ $J(\text{Red}) = \{S_1, S_2, S_3, S_4\};$ $J(\text{Blue}) = \{S_1, S_2, S_3, S_4, S_7, S_8\};$ $J(\text{White}) = \{S_2, S_3, S_4, S_6, S_8\}.$ Thus $(F, A) \cup_{\varepsilon} ((G, B) \cup_{\mathcal{R}} (H, C)) \neq ((F, A) \cup_{\varepsilon} (G, B)) \cup_{\mathcal{R}} ((F, A) \cup_{\varepsilon} (H, C)).$ Now, $K(\text{Red}) = \{S_1, S_2, S_3, S_4\};$ $K(\text{Blue}) = \{S_4\};$ $K(Yellow) = \{S_4, S_5, S_7\};$ $L(\text{Red}) = \{S_1, S_2, S_3, S_4\};$ $L(\text{Blue}) = \{S_4\};$ $L(\text{White}) = \{S_2, S_4\}.$ Thus $(F, A) \cap_{\varepsilon} ((G, B) \cap_{\mathcal{R}} (H, C)) \neq ((F, A) \cap_{\varepsilon} (G, B)) \cap_{\mathcal{R}} ((F, A) \cap_{\varepsilon} (H, C)).$ Again, we see that $I(\text{Green}) = \{S_3, S_4, S_5, S_6\};$ $I(Y$ ellow) = { S_4, S_5, S_6, S_7, S_8 }; $J(\text{Green}) = \{S_3, S_4, S_5, S_6, S_8\};$ $J(Y$ ellow) = { S_4, S_5, S_6, S_7, S_8 }; $K(\text{Green}) = \{S_3, S_4, S_5, S_6\};$ $K(\text{White}) = \{S_2, S_3, S_4\}.$ $L(\text{Green}) = \{S_4, S_5, S_6\};$ $L(Yellow) = \{S_4, S_5, S_7\};$ $M(\text{Red}) = \{S_1, S_2, S_3, S_4\};$ $M(\text{Blue}) = \{S_1, S_2, S_3, S_4, S_7\};$ $M(\text{Black}) = \{S_1, S_2, S_4, S_7\};$ $M(\text{Pink}) = \{S_2, S_3, S_5, S_7\}$ and $N(\text{Red}) = \{S_1, S_2, S_3, S_4\};$ $N(\text{Blue}) = \{S_1, S_2, S_3, S_4, S_7\};$ $N(\text{Black}) = \{S_1, S_2, S_4, S_7\};$ $N(\text{Pink}) = \{S_2, S_3, S_5, S_7\}.$ Thus $(F, A) \cup_{\varepsilon} ((G, B) \cap_{\varepsilon} (H, C)) \neq ((F, A) \cup_{\varepsilon} (G, B)) \cap_{\varepsilon} ((F, A) \cup_{\varepsilon} (H, C)).$ $M(\text{Green}) = \{S_3, S_4, S_5, S_6, S_8\};$ $M(Y$ ellow) = { S_4, S_5, S_7 }; $M(\text{White}) = \{S_2, S_3, S_4, S_6, S_8\};$ $N(\text{Green}) = \{S_3, S_4, S_5, S_6\};$ $N(Yellow) = \{S_4, S_5, S_7\};$ $N(\text{White}) = \{S_2, S_3, S_4\};$

.i

Now,

\$

بيج

 $O(\text{Red}) = \{S_1, S_2, S_3, S_4\};$ $O(\text{Blue}) = \{S_1, S_2, S_4, S_7\};$ $O(\text{Black}) = \{S_1, S_2, S_4, S_7\};$ $O(\text{Pink}) = \{S_2, S_3, S_5, S_7\}$ and $P(\text{Red}) = \{S_1, S_2, S_3, S_4\};$ $P(\text{Blue}) = \{S_1, S_2, S_4, S_7\};$ $P(\text{Black}) = \{S_1, S_2, S_4, S_7\};$ $P(\text{Pink}) = \{S_2, S_3, S_5, S_7\}.$ Thus $(F, A) \cap_{\varepsilon} ((G, B) \cup_{\varepsilon} (H, C)) \neq ((F, A) \cap_{\varepsilon} (G, B)) \cup_{\varepsilon} ((F, A) \cap_{\varepsilon} (H, C)).$ $O(\text{Greén}) = \{S_4, S_5, S_6\};$ $O(Yellow) = \{S_4, S_5, S_6, S_7, S_8\};$ $O(\text{White}) = \{S_2, S_4\};$ $P(\text{Green}) = \{S_3, S_4, S_5, S_6\};$ $P(Yellow) = \{S_4, S_5, S_6, S_7, S_8\};$ $P(\text{White}) = \{S_2, S_3, S_4\};$

1.3.18 Definition[15] **i**

Let U be an initial universal, $P(U)$ be the power set of U, E be the set of all parameter and $A, B \subseteq E$,

Let (F, A) and (G, B) be the two soft sets over a common universe U.

Then the basic intersection of the two soft sets (F, A) and (G, B) is define as the soft set

 $(H, C) = (F, A) \wedge (G, B)$ where $C = A \times B$ such that $H(e_1, e_2) = F(e_1) \cap G(e_2) \ \forall \ (e_1, e_2) \in A \times B.$

1.3.19 Definition [15]

Let U be an initial universal, $P(U)$ be the power set of U, E be the set of all parameters and $A, B \subseteq E$

Let (F, A) and (G, B) be the two soft sets over a common universe U.

Then the basic Union of the two soft sets (F, A) and (G, B) is defined as the soft sets

 $(H, C) = (F, A) \vee (G, B)$ where $C = A \times B$ such that $H(e_1, e_2) = F(e_1) U G(e_2) \ \forall \ (e_1, e_2) \in A \times B.$

1.3.20 Theorem [2]

If $(F, A), (G, B)$ and (H, C) are three soft sets over U, then

1. $((F, A) \wedge (G, B)) \wedge (H, C) = (F, A) \wedge ((G, B))$ $($ H,C)) $)$

 \mathfrak{S}

Ê

- 2. $((F, A) \vee (G, B)) \vee (H, C) = (F, A) \vee ((G, B) \vee (H, C))$
- 3. $(F, A) \wedge ((G, B) \vee (H, C)) = ((F, A) \wedge (G, B)) \vee ((F, A) \wedge (H, C))$
- 4. $(F, A) \vee ((G, B) \wedge (H, C)) = ((F, A) \vee (G, B)) \wedge ((F, A) \vee (H, C))$ ^ ^

The following remark shows that the parameter sets on both sides of the above assertions 3 and'4 are inconsistent in general.

L.8.2L Remark [2]

Let $(F, A), (G, B)$ and (H, C) be soft sets over a common universe U. The soft set $(F, A) \wedge ((G, B) \vee (H, C))$ on left side of 3 has the parameter set $A \times (B \times C)$ and the soft set.($(F, A) \wedge (G, B)$) \vee ($(F, A) \wedge (H, C)$) on right side of 3 has a set of parameters as $(A \times B) \times (A \times C)$. But in [15] we can not find any notion which ensure

 $A \times (B \times C) = (A \times B) \times (A \times C)$. Hence in Proposition 2.6 [15], two statements

1. $(F, A) \wedge ((G, B)^{^{\dagger}} \vee (H,C)) = ((F, A))$ $\mathcal{L}(\mathcal{L},\mathcal{L})$ \vee $((F, A)$ $($ $)$ $)$

2.
$$
(F, A) \vee ((G, B) \wedge (H, C)) = ((F, A) \vee (G, B)) \wedge ((F, A) \vee (H, C))
$$

are not true.

È

Chapter 2

 $\mathbf{\hat{S}}$

 $|\hat{\mathbf{r}}|$

Soft Matrix Theory and its Decision Making

In this chapter we review the paper of Naim Cagman and Serdar Engino'slu [5].

2.L Soft Matrices

$2.1.1$ Definition $[5]$

Let U be an initial universal, $P(U)$ be the power set of U, E be the set of all parameter and $A \subseteq E$.

A soft set (f_A, E) over U is defined by the set of order pairs.

$$
(f_A, E) = \{(f_A(e), e) : f_A(e) \in P(U), e \in E\}
$$

where $f_A: E \to P(U)$ such that $f_A(e) = \phi$ if e $\notin A$.

Here f_A is called approximation function of the soft set (f_A, E) . The set (f_A, E) is called e-approximate soft set. The element $f(e)$ is called the e-approximate value, which consists of related object of the perameter $e \in E$

2.L.2 Definition [5]

Let (f_A, E) be an approximate soft set over U. Then a unique subset of $U \times E$ is defined by

$$
R_A = \{(u, e) : u \in f_A(e), e \in E\}
$$

is called approximate relation.

$2.1.3$ Definition [5]

Let us define a mapping $\chi_{R_A}:U\times E\rightarrow \{0,\,1\}$ such that

$$
\chi_{R_A}(u,e) = \left\{ \begin{array}{ll} 1, & \text{if} \quad (u,e) \in R_A \\ 0, & \text{if} \quad (u,e) \notin R_A. \end{array} \right.
$$

If $U = \{u_1, u_2, u_3, \ldots, u_m\}$ and $E = \{e_1, e_2, e_3, \ldots, e_n\}$ and $A \subseteq E$, then R_A can be presented by a table as in the following form

if $a_{ij} = \chi_{R_A}(u_i, e_j)$ we define a matrix

Which is called an $m \times n$ soft matrix of the soft set (f_A, E) on a universe U. Various types of products for the elements of $SM_{m \times n}$ are defined in the following we reconsider these products.

According to the definition, soft set (f_A, E) is uniquely characterized by the matrix [a_{ij}]. It means that a soft set (f_A, E) is formally equal to its soft matrix $[a_{ij}]_{m \times n}$. Therefore we shell identify any soft set with its soft matrix and is use these two concept as interchangeable

The set of all $m \times n$ soft matrices over U will be denoted by $SM_{m \times n}$. From now on we shell delete the subscripts $m \times n$ of $[a_{ij}]_{m \times n}$ we use $[a_{ij}]$ instead of $[a_{ij}]_{m \times n}$

Example [5] $2.1.4$

 $\tilde{\epsilon}$

Assume that $U = \{u_1, u_2, u_3, u_4, u_5\}$ is a universal set and $E = \{e_1, e_2, e_3, e_4, e_5\}$ is a set of parameters. If $A = \{e_2, e_3, e_4\}$ and $f_A(e_2) = \{u_2, u_4\}, f_A(e_3) = \phi, f_A(e_4) = U$, then we write a soft set $(f_A, E) = \{(\{u_2, u_4\}, e_2), (U, e_4)\}\$ and then the relation form of (f_A, E) is written by $R_A = \{(u_2, e_2), (u_4, e_2), (u_1, e_4), (u_2, e_4), (u_3, e_4), (u_4, e_4)\}\$ hence the soft matrix $[a_{ij}]$ is written by

$$
[a_{ij}] = \left[\begin{array}{ccccc} 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & \frac{1}{2} & 0 & 1 \\ 0 & 0 & 0 & 1 \end{array}\right]
$$

 \mathbf{r}

 $\hat{\mathbf{z}}$

2.1.5 Definition [5]

Let $[a_{ij}] \in SM_{m \times n}$. Then $[a_{ij}]$ is called

- 1. A zero matrix is denoted by [0], if $a_{ij} = 0$ for all i and j
- 2. An A-universal soft matrix $[\tilde{a}_{ij}]$, if $a_{ij} = 1$ for $j \in I_A = \{j : e_j \in A\}$ and $i = 1$, $2, 3, ..., m$
- 3. A universal soft set matrix denoted by [1], if $a_{ij} = 1$ for all i and j

$2.1.6$ Example [5]

Assume that $U = \{u_1, u_2, u_3, u_4, u_5\}$ is a universal set and $E = \{e_1, e_2, e_3, e_4\}$ is a set of parameters and $[a_{ij}]$, $[c_{ij}]$, $[d_{ij}] \in SM_{5\times 4}$.

If $A = \{e_1, e_3\}$ and $f_A(e_1) = \phi$, $f_A(e_3) = \phi$ then $[a_{ij}] = [0]$ is a zero soft matrix written by

$$
[a_{ij}] = [0] = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}.
$$

If $C = \{e_1, e_2\}$ and $f_C(e_1) = U$, $f_C(e_2) = U$. Then $[c_{ij}] = [\xi_{ij}]$ is a C-Universal soft matrix written by

If $D = E$ and $f_D(e_i) = U$, for all $e_i \in D$, then $[d_{ij}] = [1]$ is a Universal soft matrix written by

r,

2.L.7 Definition [5]

Let $[a_{ij}] \in SM_{m \times n}$. Then

5\ E

- 1. $[a_{ij}]$ is the soft submatrix of $[b_{ij}]$, denoted by $[a_{ij}] \subseteq [b_{ij}]$, if $a_{ij} \leq b_{ij}$ for all i and j.
- 2. $[a_{ij}]$ is the proper soft submatrix of $[b_{ij}]$, denoted by $[a_{ij}] \subset [b_{ij}]$, if $a_{ij} \leq b_{ij}$ for at least one item $a_{ij} < b_{ij}$ all i and j.
- 3. $[a_{ij}]$ is the soft equal matrix of $[b_{ij}]$, denoted by $[a_{ij}] = [b_{ij}]$, if $a_{ij} = b_{ij}$ for all i and j.

2.1.8 Definition [5]

Let $[a_{ij}], [b_{ij}] \in SM_{m \times n}$. Then the soft matrix $[c_{ij}]$ is called

- 1. Union of $[a_{ij}]$ and $[b_{ij}]$, denoted by $[a_{ij}] \cup [b_{ij}] = [c_{ij}]$, if $[c_{ij}] = \max\{a_{ij}, b_{ij}\}$ for all i and j .
- 2. Intersection of $[a_{ij}]$ and $[b_{ij}]$, denoted by $[a_{ij}] \cap [b_{ij}] = [c_{ij}]$, if $[c_{ij}] = \min\{a_{ij}, b_{ij}\}$ for all i and j .
- 3. Complement of $[a_{ij}]$, denoted by $[a_{ij}]^{\circ} = [c_{ij}]$, if $c_{ij} = 1 a_{ij}$ for all i and j.

2.L.9 Definition [5]

Let $[a_{ij}]$, $[b_{ij}] \in SM_{m \times n}$. Then $[a_{ij}]$ and $[b_{ij}]$ are disjoint, if $[a_{ij}] \cap [b_{ij}] = [0]$ for all i and j .

2.1.10 Example [5]

Assume that
$$
[a_{ij}] = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}
$$
, $[b_{ij}] = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{bmatrix}$.
\nThen
\n $[a_{ij}]\dot{\cup}[b_{ij}] = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \end{bmatrix}$, $[a_{ij}]\dot{\cap}[b_{ij}] = [0]$, $[a_{ij}]^o = \begin{bmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix}$

I

I J

l J I

I I

2.1.11 Proposition [5]

Let $[a_{ij}] \in SM_{m \times n}$. Then

1.
$$
[[a_{ij}]^{\circ}]^{\circ} = [a_{ij}]
$$

2. $[0]$ ^o = [1].

a

r9

 \blacktriangledown

2.1.12 Proposition [5]

Let $[a_{ij}], [b_{ij}], [c_{ij}] \in SM_{m \times n}$. Then

1. $[a_{ij}] \subseteq [1]$

2.
$$
[0] \subseteq [a_{ij}]
$$

- 3. $[a_{ij}] \subseteq [a_{ij}]$
- 4. $[a_{ij}] \subseteq [b_{ij}]$ and $[b_{ij}] \subseteq [c_{ij}] \Longrightarrow [a_{ij}] \subseteq [c_{ij}].$

2.L.13 Proposition [5]

Let $[a_{ij}], [b_{ij}], [c_{ij}] \in SM_{m \times n}$. Then

1.
$$
[a_{ij}] = [b_{ij}]
$$
 and $[b_{ij}] = [c_{ij}] \Leftrightarrow [a_{ij}] = [c_{ij}]$
\n2. $[a_{ij}] \subseteq [b_{ij}]$ and $[b_{ij}] \subseteq [a_{ij}] \Leftrightarrow [a_{ij}] = [b_{ij}].$

2.L.14 Proposition [f]

Let $[a_{ij}], [b_{ij}], [c_{ij}] \in SM_{m \times n}$. Then

- 1. $[a_{ij}] \cup [a_{ij}] = [a_{ij}]$
- 2. $[a_{ij}] \cup [0] = [a_{ij}]$

3.
$$
[a_{ij}] \cup [1] = [1]
$$

4.
$$
[a_{ij}] \cup [a_{ij}]^{\circ} = [1]
$$

5.
$$
[a_{ij}]\cup [b_{ij}]=[b_{ij}]\cup [a_{ij}]
$$

6. $([a_{ij}]\cup [b_{ij}])\cup [c_{ij}] = [a_{ij}]\cup ([b_{ij}]\cup [c_{ij}]).$

2.1.15 Proposition [5]

Let $[a_{ij}], [b_{ij}], [c_{ij}] \in SM_{m \times n}$. Then

1. $[a_{ij}] \cap [a_{ij}] = [a_{ij}]$

 \mathcal{L}

انيا
الخليج

- 2. $[a_{ij}] \cap [0] = [0]$
- 3. $[a_{ij}] \cap [1] = [a_{ij}]$

4.
$$
[a_{ij}] \cap [a_{ij}]^o = [0]
$$

5.
$$
[a_{ij}] \cap [b_{ij}] = [b_{ij}] \cap [a_{ij}]
$$

6. $([a_{ij}] \cap [b_{ij}]) \cap [c_{ij}] = [a_{ij}] \cap ([b_{ij}] \cap [c_{ij}]).$

$2.1.16$ Proposition [5]

Let $[a_{ij}], [b_{ij}]$ and $[c_{ij}] \in SM_{m \times n}$. Then De Morgan's laws are valid

1.
$$
([a_{ij}] \cap [b_{ij}])^0 = [a_{ij}]^0 \cup [b_{ij}]^0
$$

2.
$$
([a_{ij}] \cup [b_{ij}])^0 = [a_{ij}]^0 \cap [b_{ij}]^0.
$$

Proof. For all i and j

$$
(\overline{a}_{ij}] \cap [b_{ij}])^0 = [\max\{a_{ij}, b_{ij}\}]^o
$$

= [1 - max\{a_{ij}, b_{ij}\}]
= [\min\{1 - a_{ij}, 1 - b_{ij}\}]
= [a_{ij}]^o \cap [b_{ij}]^o

ii.

i.

It can be proved similarly \blacksquare

2.1.17 Example [5]

Let $[a_{ij}], [b_{ij}] \in SM_{5\times 4}$ as in Example 2.1.10. Then $([a_{ij}] \cup [b_{ij}])^0 = [a_{ij}]^0 \cap [b_{ij}]^0 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 \end{bmatrix}$ and $([a_{ij}] \cap [b_{ij}])^0 = [a_{ij}]^0 \cup [b_{ij}]^0 = [1]$

2.1.18 Proposition [5]

i.
C

Let $[a_{ij}]$, $[b_{ij}]$ and $[c_{ij}] \in SM_{m \times n}$. Then

- 1. $[a_{ij}] \cup ([b_{ij}]) \cap [c_{ij}]) = ([a_{ij}] \cup [b_{ij}]) \cap ([a_{ij}] \cup [c_{ij}])$
- 2., $[a_{ij}] \cap ([b_{ij}]) \cup [c_{ij}]) = ([a_{ij}] \cap [b_{ij}]) \cap ([a_{ij}] \cap [c_{ij}]).$

2.2 Product of Soft Matrices

In this section we define four special products of soft matrices to construct soft decision making methods.

2.2.1 Definition [5]

Let $[a_{ij}], [b_{ij}] \in SM_{m \times n}$. Then And product of $[a_{ij}]$ and $[b_{ik}]$ is defined by $\wedge : SM_{m \times n} \times SM_{m \times n} \to SM_{m \times n^2},$ $[a_{ij}] \wedge [b_{ik}] = [c_{ip}]$ Where $c_{ip} = \min(a_{ij},b_{ik})$ such that $p = n(j - 1) + k$.

2.2.2 Definition.[5]

Let $[a_{ij}]$, $[b_{ij}] \in SM_{m \times n}$. Then Or- product of $[a_{ij}]$ and $[b_{ik}]$ is defined by
 $\vee : SM_{m \times n} \times SM_{m \times n} \rightarrow SM_{m \times n^2}$, $[a_{ij}] \vee [b_{ik}] = [c_{ip}]$ $\vee: SM_{m \times n} \times SM_{m \times n} \rightarrow SM_{m \times n^2},$ Where $c_{ip} = \max(a_{ij}, b_{ik})$ such that $p = n(j - 1) + k$

2.2.3 Definition [5]

Let $[a_{ij}], [b_{ij}] \in SM_{m \times n}$. Then And-Not-product of $[a_{ij}]$ and $[b_{ik}]$ is defined by
 $\overline{\wedge} : SM_{m \times n} \times SM_{m \times n} \rightarrow SM_{m \times n^2}$, $[a_{ij}] \overline{\wedge} [b_{ik}] = [c_{ip}]$ $\overline{\wedge}: SM_{m \times n} \times SM_{m \times n} \rightarrow SM_{m \times n^2},$ Where $c_{ip} = \min(a_{ij}, 1 - b_{ik})$ such that $p = n(j - 1) + k$.

2.2.4 Definition [5]

Let $[a_{ij}], [b_{ij}] \in SM_{m \times n}$. Then Or-Not- product of $[a_{ij}]$ and $[b_{ik}]$ is defined by
 $\veeeq : SM_{m \times n} \times SM_{m \times n} \rightarrow SM_{m \times n^2}$, $[a_{ij}] \veeeq [b_{ik}] = [c_{ip}]$ $Y: SM_{m \times n} \times SM_{m \times n} \rightarrow SM_{m \times n^2}$ Where $c_{ip} = \max(a_{ij}, 1 - b_{ik})$ such that $p = n(j - 1) + k$.

2.2.5 Example, [5]

Assume that $[a_{ij}]$, $[b_{ij}] \in SM_{5 \times 4}$

1 ,l I

> I I

i
_i
_international

$$
[a_{ij}] = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{bmatrix}, \qquad [b_{ij}] = \begin{bmatrix} 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

Then

$$
[a_{ij}] \wedge [b_{ik}] = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 \end{bmatrix}
$$

similarly we can find the other product $[a_{ij}] \vee [b_{ik}]$, $[a_{ij}] \wedge [b_{ik}]$, $[a_{ij}] \vee [b_{ik}]$ note that the commutativity is not valid for the soft matrices.

2.2.6 Proposition [5]

Let $[a_{ij}], [b_{ij}] \in SM_{(m \times n)}$. Then the following De Morgan's types of result are true

1.
$$
([a_{ij}] \wedge [b_{ij}])^0 = [a_{ij}]^0 \vee [b_{ij}]^0
$$

2.
$$
((a_{ij}] \vee (b_{ij}])^0 = (a_{ij}]^0 \wedge (b_{ij}]^0
$$
.

3.
$$
((a_{ij}^{\mathbf{i}})^{\vee} [b_{ij}])^{0} = [a_{ij}]^{0} \wedge [b_{ij}]^{0}
$$
.

4.
$$
([a_{ij}]\,\bar{\wedge}\,[b_{ij}])^0=[a_{ij}]^0\vee [b_{ij}]^0.
$$

2.3 Soft min-max Decision Making

In this section we construct a soft max-min decision making(SMmDM) method by using soft max-min decision function which is also defined here. The method selects optimum alternative from the set of all alternatives

2.3.L Definition [5]

Let $[c_{ij}] \in SM_{m \times n^2}$, $I_K = \{p : \exists i, c_{ip} \neq 0, (k-1)n < p \leq k\bar{n}\}$ for all $k \in I = \{1,$ 2, 3, ..., n}. Then the soft max-min decision function, denoted by Mm , is defined as follows

 $Mm:SM_{m\times n^2}\to SM_{m\times 1},$ $Mm[c_{ip}]=\left[\max_{k\in I}\{t_k\}\right]$
where where
 $t_k = \begin{cases} \min_{p \in I_k} \{c_{ip}\}, & \text{if } I_K \neq \varphi \\ 0, & \text{if } I_K = \varphi \end{cases}$

the one column soft matrix $Mm[c_{ip}]$ is called Max-min decision soft matrix.

2.3.2 Definition [5]

Let $U = \{u_1, u_2, \ldots, u_n\}$ be initial universe and $Mm[c_{ip}] = [d_{i1}]$. Then a subset of U can be obtained by using $[d_{i1}]$ as in he following way

 $opt_{[d_{i1}]}(U) = \{u_i: u_i \in U, d_{i1} = 1\}$

which is called the optimum solution.

Now, by using the definitions we can construct a SMmDM method by the following . algorithm.

Step 1: Choose feasible subsets of the set of parameters,

Step 2: construct the soft matrix for each set of parameters,

Step 3: find a convenient product of the soft matrices,

Step 4: find a max min decision soft matrix,

Step 5: find an optimum set of U.

Note that, by the similar way, we can define soft min max, soft min min and soft max max decision making methods

which may be denoted by SmMDM, SmmDM, SMMDM respectively. One of them may be more useful than others according to the type of the problems.

2.4 Applications

☜

Assume that a real estate agent has a set of different types of houses $U = \{u_1, u_2,$ u_3, u_4, u_5 which may be characterized by a set of parameters $E = \{e_1, e_2, e_3, e_4\}.$ For $j=1,2,3,4$ the parameters e_j stand for "in good location", "cheap", "modern", "large", respectively. Then we can give the following examples.

2.4.1 Example [5]

Suppose that a married couple, Mr. X and Mrs. X, come to the real estate agent to buy a house. If each partner has to consider their own set of parameters, then we select a house on the basis of the sets of partners' parameters by using the SMmDM as follows.

Assume that $U = \{u_1, u_2, u_3, u_4, u_5\}$ is a universal set and $E = \{e_1, e_2, e_3, e_4\}$ is a set of all parameters.

Step 1: First, Mr. X and Mrs. X have to choose the sets of their parameters, $A = \{e_2, e_3, e_4\}$ and $B = \{e_1, e_3, e_4\}$, respectively.

Step 2: Then we can write the following soft matrices which are constructed according to their parameters.

Step 3: Now, we can find a product of the soft matrices $[a_{ij}]$ and $[b_{ik}]$ by using And-product as follows

X and Mrs. Here, we use And-product since both Mr. X and Mrs. X's choices have to be considered.

Step 4: We can find a max min decision soft matrix as

$$
\operatorname{Mm}([a_{ij}]\wedge [b_{ik}])=\left[\begin{array}{c}1\\0\\0\\0\\0\end{array}\right]
$$

G

a s

s

Step 5: Finally, we can find an optimum set of U according to $Mm.[a_{ij}] \wedge [b_{ik}]$ $optM m._{[a_{ij}] \wedge [b_{ik}]}. (U) = \{u_1\}$

where u_1 is an optimum house to buy for Mr. X and Mrs. X.

Note that the optimal set of U may contain more than one element.

Similarly, we can also use the other products $([a_{ij}] \vee [b_{ik}])$, $([a_{ij}] \wedge [b_{ik}])$ and $([a_{ij}] \vee [b_{ik}])$ for the other convenient problems.

Chapter 3

 \in

ćΚ

ङ

Soft Matrix Theory and its **Decision Making: A New** Approach

In this chapter we are going to define new type of products which are binary and satisfies associative laws and distributive laws.

Binary-Product Of The Soft Matrices 3.1

3.1.1 Definition [5]

Let $[a_{ij}]$, $[b_{ij}] \in SM_{m \times n}$. Then,

i. And-product of $[a_{ij}]$ and $[b_{ij}]$ is defined by

$$
\wedge: SM_{m \times n} \times SM_{m \times n} \to SM_{m \times n^2}, [a_{ij}] \wedge [b_{ik}] = [c_{ip}]
$$

Where $c_{ip} = \min(a_{ij}, b_{ik})$ such that $p = n(j-1) + k$.

ii. Or-product of $[a_{ij}]$ and $[b_{ij}]$ is defined by

$$
\vee: SM_{m \times n} \times SM_{m \times n} \to SM_{m \times n^2}, [a_{ij}] \vee [b_{ik}] = [c_{ip}]
$$

Where $c_{ip} = \max(a_{ij}, b_{ik})$ such that $p = n(j - 1) + k$.

iii. And Not-product of $[a_{ij}]$ and $[b_{ij}]$ is defined by

$$
\overline{\wedge}: SM_{m \times n} \times SM_{m \times n} \to SM_{m \times n^2}, [a_{ij}] \overline{\wedge} [b_{ik}] = [c_{ip}]
$$

Where $c_{ip} = \min(a_{ij}, 1 - b_{ik})$ such that $p = n(j - 1) + k$.

iv. Or Not-product $[a_{ij}]$ and $[b_{ij}]$ is defined by

$$
\veeeq: SM_{m \times n} \times SM_{m \times n} \rightarrow SM_{m \times n^2}, [a_{ij}] \vee [b_{ik}] = [c_{ip}]
$$

Where $c_{ip} = \max(a_{ij}, 1 - b_{ik})$ such that $p = n(j - 1) + k$.

From above definition it is clear that all these products are not binary operations. Even if we consider two soft square matrices from $SM_{m \times m}$, any of above mentioned product will not give us a soft square matrix from $SM_{m \times m}$. As above mentioned products are not binary operations, therefore there is no question of associativity in these soft matrix product.

In the following products, for soft matrices are redefine so that they happen to be associate binary operations for the elements of $SM_{m \times n}$. These products will be called Binary-Product or simply we can write it as B-Products.

3.L.2 Definition

Let $[a_{ij}]$, $[b_{ij}] \in SM_{m \times n}$. Then And-B-Product of $[a_{ij}]$ and $[b_{ij}]$ is defined by

$$
\wedge \quad : \quad SM_{m \times n} \times SM_{m \times n} \to SM_{m \times n}, \ [a_{ij}] \wedge [b_{ik}] = [d_{iq}] = \begin{bmatrix} qn \\ \bigvee_{p=(q-1)n+1}^{q}(c_{ip}) \end{bmatrix}
$$

for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$.

Where $c_{ip} = \min(a_{ij}, b_{ik})$ such that $p = n(j - 1) + k$.

3.L.3 Definition

Let $[a_{ij}]$, $[b_{ij}] \in SM_{m \times n}$. Then Or-B- product of $[a_{ij}]$ and $[b_{ik}]$ is defined by

$$
\vee \quad : \quad SM_{m \times n} \times SM_{m \times n} \to SM_{m \times n}, \ [a_{ij}] \vee [b_{ik}] = [d_{iq}] = \left[\bigwedge_{p=(q-1)n+1}^{qn} (c_{ip}) \right]
$$

for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$.

Where $c_{ip} = \max(a_{ij}, b_{ik})$ such that $p = n(j - 1) + k$

3.L.4 Definition

Let $[a_{ij}], [b_{ij}] \in SM_{m \times n}$. Then And -Not-B-product of $[a_{ij}]$ and $[b_{ik}]$ is defined by

$$
\overline{\wedge} \quad : SM_{m \times n} \times SM_{m \times n} \to SM_{m \times n}, \ [a_{ij}] \overline{\wedge} [b_{ik}] = [d_{iq}] = \begin{bmatrix} qn \\ \bigvee_{p=(q-1)n+1}^{q}(c_{ip}) \end{bmatrix}
$$

for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$.

Where $c_{ip} = \min(a_{ij}, 1 - b_{ik})$ such that $p = n(j - 1) + k$

3.1.5 Definition

Let $[a_{ij}], [b_{ij}] \in SM_{m \times n}$. Then Or-Not-B-product of $[a_{ij}]$ and $[b_{ik}]$ is defined by

$$
\underline{\vee} : SM_{m \times n} \times SM_{m \times n} \to SM_{m \times n}, [a_{ij}] \underline{\vee} [b_{ik}] = [d_{iq}] = \left[\bigwedge_{p=(q-1)n+1}^{qn} (c_{ip}) \right]
$$

for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$.

Where $c_{ip} = \max(a_{ij}, 1 - b_{ik})$ such that $p = n(j - 1) + k$.

3.1.6 Theorem.

The And -B-Product is a binary product.

Proof. Let $[a_{ij}]$ and $[b_{ij}] \in SM_{m \times n}$. Then And-B-Product of $[a_{ij}]$ and $[b_{ij}]$ is defined by

 $\wedge : SM_{m \times n} \times SM_{m \times n} \rightarrow SM_{m \times n}$

$$
[a_{ij}]\wedge [b_{ik}] = [d_{iq}]
$$

Where $d_{iq} = \begin{pmatrix} qn \\ \bigvee_{p=(q-1)n+1} (e_{ip}) \end{pmatrix}$ $..., n$ for all $i = 1, 2, ..., m$ and $q = 1, 2$, where $e_{ip} = \min(a_{ij}, b_{ik})$ such that $p = n(j - 1) + k$, then $[a_{ij}] \wedge [b_{ik}] = [d_{iq}]$

3,1.7 Theorem

The Or-B-Product is a binary product.

Proof. Let $[a_{ij}]$ and $[b_{ij}] \in SM_{m \times n}$. Then Or-B-Product of $[a_{ij}]$ and $[b_{ij}]$ is defined by

$$
\vee: SM_{m \times n} \times SM_{m \times n} \to SM_{m \times n}
$$

\n
$$
[a_{ij}] \vee [b_{ik}] = [g_{i\dot{q}}]
$$

\nWhere $g_{iq} = \left(\bigwedge_{p=(q-1)n+1}^{qn} (f_{ip})\right)$ for all $i =$
\n..., n

 $f = 1, 2, ..., m$ and $q = 1, 2,$

and $f_{ip} = \max(a_{ij}, b_{ik})$ such that $p = n(j - 1) + k$

3.1.8 Theorem

The And-Not-B-Product is a binary product.

Proof. Let $[a_{ij}]$ and $[b_{ij}] \in SM_{m \times n}$. Then And-Not-B-Product of $[a_{ij}]$ and $[b_{ij}]$ is defined by

 $\bar{\wedge}:SM_{m\times n}\times SM_{m\times n}\rightarrow SM_{m\times n}$

$$
[a_{ij}] \wedge [b_{ik}] = [d_{iq}]
$$

Where $d_{iq} = \begin{pmatrix} qn \\ \bigvee_{p=(q-1)n+1} (e_{ip}) \end{pmatrix}$ for all $i = 1, 2, ..., m$ and $q = 1, 2,$
n

and $e_{ip} = \min(a_{ij}, 1 - b_{ik})$ such that $p = n(j - 1) + k$

3.1.9 Theorem

The Or-Not-B-Product is a binary product.

Proof. Let $[a_{ij}]$ and $[b_{ij}] \in SM_{m \times n}$. Then Or-Not-B-Product of $[a_{ij}]$ and $[b_{ij}]$ is defined by

 $\veeeq:SM_{m\times n}\times SM_{m\times n}\rightarrow SM_{m\times n}$

$$
[a_{ij}] \vee [b_{ik}] = [g_{iq}]
$$

where
$$
g_{iq} = \begin{pmatrix} qn \\ \bigwedge \cdot (f_{ip}) \end{pmatrix}
$$

for all $i = 1, 2, ..., m$ and $q = 1, 2, ...,$

 \boldsymbol{n}

€

÷

and $f_{ip} = \max(a_{ij}, 1 - b_{ik})$ such that $p = n(j - 1) + k$

3.1.10 Example

$(And-B-Product)$

Let $U = \{u_1, u_2, u_3\}$ be the universal set, $E = \{e_1, e_2, e_3, e_4\}$ be the set of parameters and $A = \{e_1, e_2\}, B = \{e_3, e_4\}$ be the subsets of E.

Let $f_A: E \to P(U)$ be such that $f_A(e_1) = \{u_1, u_2\}$ $f_A(e_2) = \{u_2, u_3\}$ $f_A(e_3) = f_A(e_4) = \phi$ $R_A = \{(u_1, e_1), (u_2, e_2), (u_2, e_2), (u_3, e_2)\}\$

$$
A = [a_{ij}] = \left[\begin{array}{rrr} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array} \right]
$$

and $f_B: E \to P(U)$ be such that $f_B(e_3)=U$ $f_B(e_4) = \{u_1, u_3\}$ $f_B(e_1) = f_B(e_2) = \phi$ $R_B = \{(u_1, e_3), (\u_2, e_3), (u_3, e_3), (u_1, e_4), (u_3, e_4)\}\$

$$
B = [b_{ij}] = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

$$
[a_{ij}] \wedge [b_{ik}] = \left[\begin{array}{cccc} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array}\right] \wedge \left[\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{array}\right] ,
$$

$$
[d_{ip}] = \left[\begin{array}{cccccccccccc} 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{array}\right]
$$

and $d_{ip} = \min(a_{ij}, b_{ik})$ such that $p = n(j-1) + k$

Where
$$
y_{iq} = \begin{pmatrix} q^4 \\ \bigvee_{p=(q-1)4+1}^{q} (d_{ip}) \end{pmatrix}
$$

for all
$$
i = 1, 2, 3
$$
 and $q = 1, 2, 3, 4$

Then

 \widehat{S}

惫

 $[Y] = [y_{iq}]_{3 \times 4} = [a_{ij}] \wedge [b_{ik}]$

$$
[a_{ij}] \wedge [b_{ik}] = \left[\begin{array}{cccc} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array}\right]
$$

$3.1.11$ Example

(Or-B-Product)

Let $U = \{u_1, u_2, u_3\}$ be the universal set, $E = \{e_1, e_2, e_3, e_4\}$ be the set of parameters and $A = \{e_1, e_2\}, B = \{e_3, e_4\}$ be the subsets of E.

Let $f_A: E \to P(U)$, be such that

 $f_A(e_1) = \{u_1, u_2\}$ $f_{A}^{*}(e_{2}) = \{u_{2}, u_{3}\}\$ $f_A(e_3) = f_A(e_4) = \phi$ $R_A = \{(u_1, e_1), (u_2, e_2), (u_2, e_2), (u_3, e_2)\}\$

and $f_B: E \to P(U)$ be such that $f_B(e_3)=U$ $f_B(e_4) = \{u_1, u_3\}$ $f_B(e_1) = f_B(e_2) = \phi$ $R_B = \{(u_1, e_3), (u_2, e_3), (u_3, e_3), (u_1, e_4), (u_3, e_4)\}\$

$$
B = [b_{ij}] = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

$$
[a_{ij}] \vee [b_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \vee \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

3. Soft Matrix Theory and its Decision Making: A New Approach

and f_{ip} = max(a_{ij} , b_{ik}) such that $p = n(j - 1) + k$

if
$$
x_{iq} = \left(\bigwedge_{p=(q-1)4+1}^{4} (f_{ip})\right)
$$

for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

_{SO}

E

$$
[X] = [x_{iq}]_{3\times 4} = [a_{ij}] \vee [b_{ik}]
$$

$$
[a_{ij}] \vee [\check{b}_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

3.1.12 **Example**

(And-Not-B-Product)

Let $U = \{u_1, u_2, u_3\}$ be the universal set, $E = \{e_1, e_2, e_3, e_4\}$ be the set of parameters and $A = \{e_1, e_2\}, B = \{e_3, e_4\}$ be the subsets of E.

Let $f_A: E \to P(U)$ be such that $f_A(e_1) = \{u_1, u_2\}$ $f_A(e_2) = \{u_2, u_3\}$ $f_A(e_3) = f_A(e_4) = \phi$ $R_A = \{(u_1, e_1), (u_2, e_2), (u_2, e_2), (u_3, e_2)\}\$

and $f_B: E \to P(U)$ be such that $f_B(e_3)=U$ $f_B(e_4) = \{u_1, u_3\}$

$$
f_B(e_1) = f_B(e_2) = \phi
$$

\n
$$
R_B = \{(u_1, e_3), (u_2, e_3), (u_3, e_3), (u_1, e_4), (u_3, e_4)\}\
$$

$$
B = [b_{ij}] = \left[\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{array}\right]
$$

 $[a_{ij}]\barwedge [b_{ik}] = \left[\begin{array}{cccc} 1 & 0 & 0 & 0\ 1 & 1 & 0 & 0\ 0 & 1 & 0 & 0 \end{array}\right] \barwedge \left[\begin{array}{cccc} 0 & 0 & 1 & 1\ 0 & 0 & 1 & 0\ 0 & 0 & 1 & 1 \end{array}\right] \quad ^*$

and $d_{ip} = \min(a_{ij}, 1 - b_{ik})$ such that $p = n(j - 1) + k$

where
$$
y_{iq} = \begin{pmatrix} q^4 \\ \bigvee_{p=(q-1)4+1} (d_{ip}) \end{pmatrix}
$$

for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

Then

۳

 \mathbf{C}

 $[Y] = [y_{iq}]_{3 \times 4} = [a_{ij}] \bar{w} [b_{ik}]$

$$
[a_{ij}]\barwedge [b_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

3.1.13 Example

(Or-Not-B-Product)

Let $U = \{u_1, u_2, u_3\}$ be the universal set, $E = \{e_1, e_2, e_3, e_4\}$ be the set of parameters and $A = \{e_1, e_2\}, B = \{e_3, e_4\}$ be the subsets of E.

Let $f_A: E \to P(U)$ be such that

ŧ

 $f_A(e_1) = \{u_1, u_2\}$ $f_A(e_2) = \{u_2, u_3\}$ $f_A(e_3) = f_A(e_4) = \phi$ $R_A = \{(u_1, e_1), (u_2, e_2), (u_2, e_2), (u_3, e_2)\}\$

۳

and $f_B: E \to P(U)$ be such that $f_B(e_3) = U$ $f_B(e_4) = \{u_1, u_3\}$ $f_B(e_1) = f_B(e_2) = \phi$ $R_B = \{(u_1, e_3), (u_2, e_3), (u_3, e_3), (u_1, e_4), (u_3, e_4)\}\$

 $\overline{}$

$$
B = [b_{ij}] = \left[\begin{array}{rrr} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{array} \right]
$$

$$
[a_{ij}] \vee [b_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \vee \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

$$
[f_{ip}] = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 \end{bmatrix}
$$

and $f_{ip} = \max(a_{ij}, 1-b_{ik})$ such that $p = n(j-1) + k$

Where
$$
x_{iq} = \left(\bigwedge_{p=(q-1)4+1}^{4} (f_{ip})\right)
$$

for all
$$
i = 1, 2, 3
$$
 and $q = 1, 2, 3, 4$

29

{SO} $[X] = [x{iq}]_{n \times d} = [a_{ij}] \vee [b_{ik}]$

$$
[a_{ij}] \vee [b_{ik}] = \left[\begin{array}{cccc} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array}\right]
$$

so the above example shows that the defined product is a binary operation

3.1.14 Theorem

 $71 + 16063$

The associative law holds with respect to And-B-Product.

Proof. Let $[a_{ij}]$, $[b_{ij}]$, $[c_{ij}] \in SM_{m \times n}$. Then And-B-Product of $[a_{ij}]$ and $[b_{ij}]$ is define by $\wedge: SM_{m \times n} \times SM_{m \times n} \rightarrow SM_{m \times n}$

$$
[a_{ij}] \wedge [b_{ik}] = [d_{iq}]
$$

where
$$
d_{iq} = \begin{pmatrix} qn \\ \bigvee_{p=(q-1)n+1} (e_{ip}) \end{pmatrix}
$$

for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$

and $e_{ip} = \min(a_{ij}, b_{ik})$ such that $p = n(j-1) + k$ now

 $([a_{ij}] \wedge [b_{ik}]) \wedge [c_{ij}] = [d_{ig}] \wedge [c_{ij}]$

$$
([a_{ij}]\wedge [b_{ik}])\wedge [c_{ij}]=[h_{iq}]
$$

where $h_{iq} = \left(\bigvee_{n=(a-1)n+1}^{qn}(s_{ip})\right)$

for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$

and $s_{ip} = \min(d_{ij}, c_{ik})$ such that $p = n(j - 1) + k$ now R.H.S

$$
[b_{ij}] \wedge [c_{ik}] = [g_{iq}]
$$

Where $g_{iq} = \left(\bigvee_{n=(q-1)n+1}^{qn} (f_{ip}) \right)$

for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$

and $f_{ip} = \min(b_{ij}, c_{ik})$ such that $p = n(j - 1) + k$ $[a_{ij}] \wedge ([b_{ik}] \wedge [c_{ij}]) = [a_{ij}] \wedge ([g_{iq}])$

3. Soft Matrix Theory and its Decision Making: A New Approach

Now

$$
[a_{ij}] \wedge ([b_{ik}] \wedge [c_{ij}]) = [h_{iq}]
$$

Where $h_{iq} = \begin{pmatrix} qn \\ \sqrt{n} & (t_{ip}) \end{pmatrix}$ for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$
and $t_{ip} = \min(a_{ij}, g_{ik})$ such that $p = n(j - 1) + k$
than

$$
([a_{ij}] \wedge [b_{ik}]) \wedge [c_{ij}] = [a_{ij}] \wedge ([b_{ik}] \wedge [c_{ij}]).
$$

3.1.15 Theorem

 $|\mathcal{B}\rangle$

The associative law holds with respect to Or-B-Product.

Proof. Let $[a_{ij}]$, $[b_{ij}]$, $[c_{ij}] \in SM_{m \times n}$. Then Or-Product of $[a_{ij}]$ and $[b_{ij}]$ is define by $\vee: SM_{m \times n} \times SM_{m \times n} \rightarrow SM_{m \times n}$

$$
[a_{ij}] \vee [b_{ik}] = [d_{iq}]
$$

Where
$$
d_{iq} = \left(\bigwedge_{p=(q-1)n+1}^{qn} (e_{ip})\right)
$$
 for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$

and $e_{ip} = \max(a_{ij}, b_{ik})$ such that $p = n(j-1) + k$ $(a_{ij} \vee [b_{ik}] \vee [c_{ij}] = [d_{iq}] \vee [c_{ij}]$

$$
([a_{ij}] \vee [b_{ik}]) \vee [c_{ij}] = [h_{iq}]
$$

Where $h_{iq} = \left[\bigwedge_{n=(n-1)n+1}^{qn} (s_{ip}) \right]$

for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$

and $s_{ip} = \max(d_{ij}, c_{ik})$ such that $p = n(j - 1) + k([a_{ij}] \vee [b_{ik}]) \vee [c_{ij}] = [h_{iq}]$ now R.H.S

$$
[b_{ij}]\vee[c_{ik}]=[g_{iq}]
$$

Where $g_{iq} = \left(\bigwedge_{p=(q-1)n+1}^{qn} (f_{ip})\right)$ 'for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$

and $f_{ip} = \max(b_{ij}, c_{ik})$ such that $p = n(j - 1) + k$ \sim Now

 $[a_{ij}] \vee ([b_{ik}] \vee [c_{ij}]) = [a_{ij}] \vee ([g_{iq}])$

$$
[a_{ij}] \vee ([b_{ik}] \vee [c_{ij}]) = [h_{iq}]
$$

Where
$$
h_{iq} = \left(\bigwedge_{p=(q-1)n+1}^{qn}(t_{ip})\right)
$$
 for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$ and $t_{ip} = \max(a_{ij}, g_{ik})$ such that $p = n(j-1) + k$ $([a_{ij}] \vee [b_{ik}]) \vee [c_{ij}] = [a_{ij}] \vee ([b_{ik}] \vee [c_{ij}]).$

3.1.16 Theorem

Ar ℓ

 $|\mathcal{B}\rangle$

The associative law holds with respect to And-Not-B-Product.

 $([a_{ij}]\,\overline{\wedge}\,[b_{ik}])\,\overline{\wedge}\,[c_{ij}]=[a_{ij}]\,\overline{\wedge}\,([b_{ik}]\,\overline{\wedge}\,[c_{ij}])$ Proof. Strightforword. ■

3.L.17 Theorem

The associative law holds with respect to Or-Not-B-Product.

 $([a_{ij}] \vee [b_{ik}]) \vee [c_{ij}] = [a_{ij}] \vee [b_{ik}] \vee [c_{ij}])$ Proof. Strightforword. ■

3.1.18 Example

(Associative law with respect to And-B-Product)

Let $U = \{u_1, u_2, u_3\}$ be the universal set, $E = \{e_1, e_2, e_3, e_4\}$ be the set of parameters and $A = \{e_1, e_2\}, B = \{e_3, e_4\}, C = \{e_2, e_3\}$ be the subsets of E.

Let $f_A: E \to P(U)$ be such that $f_A(e_1) = \{u_1, u_2\}$ $f_A(e_2) = \{u_2, u_3\}$

 $f_A(e_3) = f_A(e_4) = \phi$

$$
R_A = \{(u_1, e_1), (u_2, e_2), (u_2, e_2), (u_3, e_2)\}
$$

Ý.

and $f_B : E \to P(U)$ be such that $f_B(e_3) = U$ $f_B(e_4) = \{u_1, u_3\}$

 $f_B(e_1) = f_B(e_2) = \phi$ $R_B = \{(u_1, e_3), (u_2, e_3), (u_3, e_3), (u_1, e_4), (u_3, e_4)\}\$

and $f_C: E \to P(U)$ be such that $f_C(e_2) = \{u_2\}$ $f_C(e_3) = \{u_2, u_3\}$ $f_C(e_1) = f_C(e_4) = \phi$ $R_C = \{(u_2, e_2), (u_2, e_3), (u_3, e_3)\}\$

Now to prove $([a_{ij}] \wedge [b_{ik}]) \wedge [c_{ij}] = [a_{ij}] \wedge ([b_{ik}] \wedge [c_{ij}])$ Firstly we Find that $([a_{ij}] \wedge [b_{ik}]) \wedge [c_{ij}]$

$$
[a_{ij}] \wedge [b_{ik}] = \left[\begin{array}{cccc} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array}\right] \wedge \left[\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{array}\right]
$$

and $d_{ip} = \min(a_{ij}, b_{ik})$ such that $p = n(j-1) + k$

Ĩ,

・やんとす

 $\Gamma_{\rm c}=\Gamma_{\rm d}$

Where
$$
y_{iq} = \begin{pmatrix} q4 \\ \bigvee_{p=(q-1)4+1} (d_{ip}) \end{pmatrix}
$$

so

for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

 $[Y]=[y_{iq}]_{3\times4}=\big[a_{ij}\big]\wedge\big[b_{ik}\big]$

$$
[y_{iq}]_{3\times 4} = [a_{ij}] \wedge [b_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

Now

T)

 $\begin{picture}(20,20) \put(0,0){\line(1,0){15}} \put(15,0){\line(1,0){15}} \put(15,0){\line(1$

 $([a_{ij}]\wedge [b_{ik}])\wedge [c_{ij}]=[{\dot{y}_{ij}}]\wedge [c_{ik}]$

$$
[y_{ij}] \wedge [c_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \wedge \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}
$$

and
$$
e_{ip} = \min(y_{ij}, c_{ik})
$$
 such that $p = n(j-1) + k$

Where $w_{iq} = \begin{pmatrix} q^4 \\ \bigvee^{q^4}_{p=(q-1)4+1} (e_{ip}) \end{pmatrix}$ for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$ **SO**

$$
[W] = [w_{iq}]_{3 \times 4} = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

 $[w_{iq}] = [y_{ij}] \wedge [c_{ik}]$ $[w_{iq}] = ([a_{ij}] \wedge [b_{ik}]) \wedge [c_{ij}]$

$$
([a_{ij}] \wedge [b_{ik}]) \wedge [c_{ij}] = \left[\begin{array}{cccc} 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array}\right]
$$

 \mathbf{r}^{\star}

Now we find $[a_{ij}] \wedge ([b_{ik}] \wedge [c_{ij}])$

$$
[b_{ik}] \wedge [c_{ij}] = \left[\begin{array}{cccc} 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{array}\right] \wedge \left[\begin{array}{cccc} 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 \end{array}\right]
$$

 ${\rm Let}$

 \mathfrak{g}

and
$$
f_{ip} = \min(b_{ij}, c_{ik})
$$
 such that $p = n(j-1) + k$

Where
$$
v_{iq} = \begin{pmatrix} q^4 \\ \sqrt{ } \\ p=(q-1)4+1 \end{pmatrix}
$$
 for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$
\nso
\n
$$
[V] = [v_{iq}]_{3\times 4} = [b_{ik}] \wedge [c_{ij}]
$$
\n
$$
[v_{iq}] = [b_{ik}] \wedge [c_{ij}] = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$
\n
$$
[a_{ij}] \wedge ([b_{ik}] \wedge [c_{ij}]) = [a_{ij}] \wedge [v_{iq}]
$$
\n
$$
[a_{ij}] \wedge [v_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \wedge \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$
\n
$$
[g_{ip}] = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}
$$
\nand $g_{ip} = \min(b_{ij}, c_{ik})$ such that $p = n(j - 1) + k$
\nWhere $s_{iq} = \begin{bmatrix} q^4 \\ \sqrt{ } & (g_{ip}) \\ p = (q-1)4+1 \end{bmatrix}$ for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

$$
[S] = [s_{iq}]_{3 \times 4} = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

3. Soft Matrix Theory and its Decision Making: A New Approach

 $[s_{iq}] = [a_{ij}] \wedge [v_{iq}]$ $[s_{iq}] = [a_{ij}] \wedge ([b_{ik}] \wedge [c_{ij}])$ SO

there fore

ê

 $([a_{ij}] \wedge [b_{ik}]) \wedge [c_{ij}] = [a_{ij}] \wedge ([b_{ik}] \wedge [c_{ij}])$ now it can satisfy the associative property

3.1.19 Example

(Associative law over Or-B-Product)

Let $U = \{u_1, u_2, u_3\}$ be the universal set, $E = \{e_1, e_2, e_3, e_4\}$ be the set of parameters and $A = \{e_1, e_2\}, B = \{e_3, e_4\}, C = \{e_2^*, e_3\}$ be the subsets of E.

Let $f_A: E \to P(U)$ be such that $f_A(e_1) = \{u_1, u_2\}$ $f_A(e_2) = \{u_2, u_3\}$ $\frac{1}{2}$ $f_A(e_3) = f_A(e_4) = \phi$. $R_A = \{(u_1, e_1), (u_2, e_2), (u_2, e_2), (u_3, e_2)\}\$ λ - $\begin{array}{c|cc} R_A & e_1 & e_2 & e_3 \\ \hline \hline \dot{u_1} & 1 & 0 & 0 \\ \hline \dot{u_2} & 1 & 1 & 0 \\ \hline \dot{u_3} & 0 & 1 & 0 \end{array}$ $\overline{0}$ $\overline{0}$. $\overline{0}$ * * * * $A = [a_{ij}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \end{bmatrix}$ $\begin{bmatrix} 0 & 1 & 0 \end{bmatrix}$ and $f_B: E \to P(U)$ be such that $f_B(e_3)=U$ \pmb{f} .

 $f_B(e_4) = \{u_1, u_3\}$

 $f_B(e_1) = f_B(e_2) = \phi$

$$
R_B = \{(u_1, e_3), (u_2, e_3), (u_3, e_3), (u_1, e_4), (u_3, e_4)\}
$$

$$
[Y] = [y_{iq}] = [a_{ij}] \vee [b_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

Now

E)

 $([a_{ij}] \vee [b_{ik}]) \vee [c_{ij}] = [y_{iq}] \vee [c_{ij}]$

$$
[y_{ij}] \vee [c_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \vee \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}
$$

$$
[\hat{e}_{ip}] = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \end{bmatrix}
$$
and $e_{ik} = \max(y_{ik}, c_{ik})$ such that $y = y(i-1) + k$

Where
$$
w_{iq} = \begin{pmatrix} q4 \\ \bigwedge_{p=(q-1)4+1}^{q4}(e_{ip}) \end{pmatrix}
$$
 for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

$$
_{\rm\bf SO}
$$

$$
[\hat{W}] = [w_{iq}]_{3 \times 4} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

 $[w_{iq}] = [y_{ij}] \vee [c_{ik}]$ $[w_{iq}] = ([a_{ij}] \vee [b_{ik}]) \vee [c_{ij}]$

$$
([a_{ij}] \vee [b_{ik}]) \vee [c_{ij}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

Now we find $[a_{ij}] \vee ([b_{ij}] \vee [c_{ij}])$

$$
[b_{ij}] \vee [c_{ik}] = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix} \vee \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}
$$

 $37\,$

 $\mathbf{1}$

Let

 $\hat{\mathbf{z}}$

 \mathbb{Z}^n .

1,..-

$$
[f_{ip}] = \left[\begin{array}{cccccccccccc} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{array} \right]
$$

and
$$
f_{ip} = \max(b_{ij}, c_{ik})
$$
 such that $p = n(j-1) + k$

Where
$$
v_{iq} = \left(\bigwedge_{p=(q-1)4+1}^{q4} (f_{ip})\right)
$$
 for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$
so

$$
\mathcal{L}_{\mathcal{A}}(x)
$$

$$
[v_{iq}] = \left[\begin{array}{cccc} 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{array} \right]
$$

 $[V] = [v_{iq}]_{3\times4} = [b_{ij}] \vee [c_{ik}]$

$$
[b_{ij}] \vee [c_{ik}] = \left[\begin{array}{cccc} 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{array} \right]
$$

now

 $[a_{ij}] \vee ([b_{ij}] \vee [c_{ij}]) = [a_{ij}] \vee [v_{iq}]$

$$
[a_{ij}] \vee [v_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \vee \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

$$
[g_{ip}] = \left[\begin{array}{cccccccccccc} 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \end{array} \right]
$$

and g_{ip} = max (b_{ij}, c_{ik}) such that $p = n(j - 1) + k$

Where
$$
s_{iq} = \begin{bmatrix} q4 \\ \bigwedge_{p=(q-1)4+1}^{q4}(g_{ip}) \end{bmatrix}
$$
 for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$
so

I

t I t

$$
[S] = [s_{iq}]_{3 \times 4} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

as $[s_{iq}] = [a_{ij}] \vee [v_{ij}]$ $[s_{iq}] = [a_{ij}] \vee ([b_{ij}] \vee [c_{ij}])$ _{SO}

$$
[a_{ij}] \vee ([b_{ij}] \vee [c_{ij}]) = \left[\begin{array}{cccc} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array}\right]
$$

there fore

 $([a_{ij}] \vee [b_{ik}]) \vee [c_{ij}] = [a_{ij}] \vee ([b_{ik}] \vee [c_{ij}])$ Similarly we can prove $((a_{ij} \nvert \nabla [b_{ik}]) \nabla [c_{ij}] = [a_{ij} \nvert \nabla ([b_{ik} \nvert \nabla [c_{ij}])]$ $([a_{ij}] \vee [b_{ik}]) \vee [c_{ij}] = [a_{ij}] \vee ([b_{ik}] \vee [c_{ij}])$

3.1.20 Theorem

i.
K

Or-B-Product is distributive over And-B-Product. **Proof.** Let $[a_{ij}]$, $[b_{ij}]$, $[c_{ij}] \in SM_{m \times n}$. Then

$$
[b_{ij}]\wedge[c_{ik}]=[d_{iq}]
$$

Where $d_{iq} = \begin{pmatrix} qn \\ \bigvee \ \psi=(q-1)n+1 \end{pmatrix} (e_{ip})$

for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$

 \boldsymbol{n}

and $e_{ip} = \min(a_{ij}, b_{ik})$ such that $p = n(j - 1) + k$ now

 $[a_{ij}] \vee ([b_{ij}] \wedge [c_{ij}]) = [a_{ij}] \vee [d_{ik}]$

$$
[a_{ij}] \vee [d_{ik}] = [g_{iq}]
$$

if
$$
g_{iq} = \begin{pmatrix} qn \\ \bigwedge_{p=(q-1)n+1}^{qn}(f_{ip}) \end{pmatrix}
$$
 for all $i = 1, 2, ..., m$ and $q = 1, 2, ...$

Where $f_{ip} = \max(a_{ij}, b_{ik})$ such that $p = n(j - 1) + k$ $[a_{ij}] \vee ([b_{ij}] \wedge [c_{ij}]) = [a_{ij}] \vee [d_{ik}] = [g_{ij}]$

Now R.H.S

 \mathbf{e}

!

 $\hat{\vec{c}}$

$$
[a_{ij}]\wedge [b_{ik}]=[h_{iq}]
$$

 $\text{Where}\,\,h_{iq}=\left(\begin{array}{cc} & qn\ &\bigwedge\quad &(t_{ip})\end{array}\right).$ $\sqrt{p=(q-1)n+1}$ / for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$

and $t_{ip} = \max(a_{ij}, b_{ik})$ such that $p = n(j-1)+k$ Now

$$
[a_{ij}]\vee[c_{ik}]=[s_{iq}]
$$

Where s_{ij} = $\binom{1}{p=(q-1)n+1}$ for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$

And $v_{ip} = \max(a_{ij}, c_{ik})$ such that $p = n(j-1) + k$

$$
[h_{ij}]\wedge[s_{ip}]=[x_{iq}]
$$

Where $x_{iq} =$ for all $i = 1, 2, ..., m$ and $q = 1, 2, ..., n$ $\bigg\downarrow_{p=(q-1)n+1}$

And $y_{ip} = \min(h_{ij}, s_{ik})$ such that $p = n(j - 1) + k$ SO $[a_{ij}] \vee (b_{ij}] \wedge [c_{ij}]) = (a_{ij}] \wedge [b_{ij}] \wedge (a_{ij}] \vee [c_{ij}] \equiv$

3.L.21 Example

(Or-B-product is distributive over And-B-Product)

Let $U = \{u_1, u_2, u_3\}$ be the universal set, $E = \{e_1, e_2, e_3, e_4\}$ be the set of parameters and, $A = \{e_1, e_2\}, B = \{e_3, e_4\}, C = \{e_2, e_3\}$ be the subsets of E.

Let $f_A: E \to P(U)$ be such that

$$
f_A(e_1) = \{u_1, u_2\}
$$

\n
$$
f_A(e_2) = \{u_2, u_3\}
$$

\n
$$
f_A(e_3) = f_A(e_4) = \phi
$$

\n
$$
R_A = \{(u_1, e_1), (u_2, e_2), (u_2, e_2), (u_3, e_2)\}
$$

and $f_B: E \to P(U)$ be such that $f_B(e_3)=U$ $f_B(e_4) = \{u_1, u_3\}$ $f_B(e_1) = f_B(e_2) = \phi$ $R_B = \{(u_1, e_3), (u_2, e_3), (u_3, e_3), (u_1, e_4), (u_3, e_4)\}\$

and $f_C: E \to P(U)$ be such that

۲

 \vec{r}

 \mathcal{L}

 $f_C(e_2) = \{u_2\}$ $f_C(e_3) = \{u_2, u_3\}$ $f_C(e_1) = f_C(e_4) = \phi$ $R_C = \{(u_2, e_2), (u_2, e_3), (\tilde{u}_3, e_3)\}\$

$$
C = [c_{ij}] = \left[\begin{array}{rrr} 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 \end{array} \right]
$$

To prove $[a_{ij}] \vee ([b_{ij}] \wedge [c_{ij}]) = ([a_{ij}] \vee [b_{ij}]) \wedge ([a_{ij}] \vee [c_{ij}])$

 $L.H.S$ firstly we find

Ç

 \mathbb{R}

$$
[b_{ij}]\wedge [c_{ij}] = \begin{bmatrix} .0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix} \wedge \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}
$$

$$
[g_{ip}] = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}
$$

and g_{ip} = min(b_{ij} , c_{ik}) such that $p = n(j - 1) + k$

Where
$$
s_{iq} = \left(\bigvee_{p=(q-1)n+1}^{q4} (g_{ip})\right)
$$
 for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

then

$$
[s_{iq}] = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}.
$$

_{SO}

$$
[s_{iq}] = [b_{ij}] \wedge [c_{ij}] = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

Now

 $[a_{ij}] \vee ([b_{ij}] \wedge [c_{ij}]) = [a_{ij}] \vee [s_{iq}]$

$$
[a_{ij}] \vee [s_{ik}] = \left[\begin{array}{cccc} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array}\right] \vee \left[\begin{array}{cccc} 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{array}\right]
$$

 $[f_{ip}] \begin{array}{c} \end{array} \hspace{0.2cm} = \hspace{0.2cm} \left[\begin{array}{cccccccccccccccc} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \end{array} \right]$

and f_{ip} = $\max(\ddot{a}_{ij}, s_{ik})$ such that $p = n(j-1) + k$

Where
$$
e_{iq} = \left(\bigwedge_{p=(q-1)4+1}^{q4} (f_{ip})\right)
$$
 for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

$$
[a_{ij}] \vee [s_{iq}] = [e_{iq}] = \left[\begin{array}{cccc} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array}\right]
$$

 \cdot Now R.H.S

麀

 \mathcal{P}

હિ

$$
([a_{ij}]\vee [b_{ij}]) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \vee \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

$$
[d_{ip}] = \left[\begin{array}{ccccccccccc} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \end{array} \right]
$$

and $d_{ip} = \max(a_{ij}, b_{ik})$ such that $p = n(j - 1) + k$

where $x_{iq} = \begin{pmatrix} q4 \ \bigwedge^{q4} (d_{ip}) \ p=(q-1)4+1 \end{pmatrix}$ for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$ $([a_{ij}]\vee [b_{ij}]) = [x_{iq}] = \left[\begin{array}{cccc} 1 & 0 & 0 & 0 \ 1 & 1 & 0 & 0 \ 0 & 1 & 0 & 0 \end{array} \right]$

Now

$$
([a_{ij}] \vee [c_{ij}]) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \vee \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix} ,
$$

$$
[y_{ip}] = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \end{bmatrix} .
$$

and y_{ip} = max (a_{ij}, c_{ik}) such that $p = n(j - 1) + k$

43

ŗ

Where
$$
z_{iq} = \left(\bigwedge_{p=(q-1)4+1}^{q4} (y_{ip})\right)
$$
 for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

$$
[z_{iq}] = ([a_{ij}] \vee [c_{ij}]) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

now

rE

Ë

$$
([x_{ij}]\wedge [z_{ik}]) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \wedge \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

$$
[g_{ip}] = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}
$$

and g_{ip} = $\min(x_{ij}, z_{ik})$ such that $p = n(j - 1) + k$

Where
$$
h_{iq} = \begin{pmatrix} q4 \\ \bigvee_{p=(q-1)n+1}^{q4}(u_{ip}) \end{pmatrix}
$$

for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

,|

 $\ddot{\mathrm{a}}$

then

$$
[h_{iq}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

\n
$$
([x_{ij}] \wedge [z_{ik}]) = [h_{iq}]
$$

\nso $([a_{ij}] \vee [b_{ij}]) \wedge ([a_{ij}] \vee [c_{ij}]) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$
\nso
\n
$$
[a_{ij}] \vee ([b_{ij}] \wedge [c_{ij}]) = ([a_{ij}] \vee [b_{ij}]) \wedge ([a_{ij}] \vee [c_{ij}]).
$$

3.L.22 Theorem

And-B-Product is distributive over Or-B-Product. $[a_{ij}] \wedge ([b_{ij}] \vee [c_{ij}]) = ([a_{ij}] \wedge [b_{ij}]) \vee ([a_{ij}] \wedge [c_{ij}])$
Proof Strichtforword = **Proof.** Strightforword \blacksquare

3.1.23' Remark

rtr

I Ξ l'

 $\sum_{i=1}^{n}$ iF'

 $\overline{}$

And-Not-B-Product is not distributive over Or-Not-B-Product. $[a_{ij}]\bar{\wedge}([b_{ij}]\vee c_{ij}])\neq ([a_{ij}]\bar{\wedge} [b_{ij}])\vee (a_{ij}]\bar{\wedge} [c_{ij}])$

3.L.24 Example

Let $U = \{u_1, u_2, u_3\}$ be the universal set, $E = \{e_1, e_2, e_3, e_4\}$ be the set of parameters and $A = \{e_1, e_2\}, B = \{e_3, e_4\}, C = \{e_2, e_3\}$ be the subsets of E.

Let $f_A : E \to P(U)$ be such that $f_A(e_1) = \{u_1, u_2\}$ $f_A(e_2) = \{u_2, u_3\}$ $f_A(e_3) = f_A(e_4) = \phi$ $R_A = \{(u_1, e_1), (u_2, e_2), (u_2, e_2), (u_3, e_2)\}\$

 $\begin{bmatrix} 0 & 1 & 0 & 0 \end{bmatrix}$

and $f_B: E \to P(U)$ be such that $f_B(e_3) = U$ $f_B(e_4) = \{u_1, u_3\}$ $f_B(e_1) = f_B(e_2) = \phi$ $R_B = \{(u_1, e_3), (u_2, e_3), (u_3, e_3), (u_1, e_4), (u_3, e_4)\}\$

$$
B = [b_{ij}] = \left[\begin{array}{cccc} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{array} \right]
$$

and $f_C : E \to P(U)$ be such that $f_C(e_2) = \{u_2\}$

 \sqrt{R}

 $f_C(e_3) = \{u_2, u_3\}$ $f_C(e_1) = f_C(e_4) = \phi$ $R_C = \{(u_2, e_2), (u_2, e_3), (u_3, e_3)\}\$

and $f_{ip} = \max(b_{ij}, 1 - c_{ik})$ such that $p = n(j - 1) + k$

Where $e_{iq} = \left(\bigwedge_{p=(q-1)4+1}^{q4} (f_{ip})\right)$ for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$ $[e_{iq}] = [b_{ij}] \vee [c_{ik}] = \left[\begin{array}{cccc} 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \end{array} \right]$

Now

 $[a_{ij}] \overline{\wedge} ([b_{ij}] \vee c_{ij}]) = [a_{ij}] \overline{\wedge} [e_{iq}]$

$$
[a_{ij}]\barwedge [e_{ik}] = \left[\begin{array}{cccc} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array}\right] \barwedge \left[\begin{array}{cccc} 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \end{array}\right]
$$

and $g_{ip} = \min(a_{ij}, 1 - e_{ik})$ such that $p = n(j - 1) + k$

Where
$$
s_{iq} = \begin{pmatrix} q^4 \\ \sqrt{(-q-1)n+1} \end{pmatrix}
$$
 for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

then

P

$$
[s_{iq}] = \left[\begin{array}{cccc} 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array}\right]
$$

_{SO}

$$
[a_{ij}] \overline{\wedge} [e_{ik}] = [s_{iq}] = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

$$
[a_{ij}] \overline{\wedge} ([b_{ij}] \vee [c_{ij}]) = [s_{iq}] = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

Now R.H.S

 $([a_{ij}]\barwedge [b_{ij}]) \vee ([a_{ij}]\barwedge [c_{ij}])$

$$
[a_{ij}]\bar{\wedge}[b_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \bar{\wedge} \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

$$
[f_{ip}] = \begin{bmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}
$$

and $f_{ip} = \min(a_{ij}, 1 - b_{ik})$ such that $p = n(j - 1) + k$

Where $x_{iq} = \begin{pmatrix} q4 \\ \bigvee_{p=(q-1)n+1}^{q4} (f_{ip}) \end{pmatrix}$

for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

 $then$

$$
[x_{iq}]=\left[\begin{array}{cccc}1&0&0&0\\1&1&0&0\\0&1&0&0\end{array}\right]
$$

48

SO

 $\frac{1}{\sqrt{2}}$

$$
[a_{ij}]\bar{\wedge}[b_{ik}] = [x_{iq}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

$$
([a_{ij}]\bar{\wedge}[c_{ik}]) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \bar{\wedge} \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}
$$

$$
[g_{ip}] = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}
$$

and
$$
q_{in}
$$
 = min(a_{ij} , 1 - c_{ik}) such that $p = n(j-1) + k$

Where
$$
y_{iq} = \begin{pmatrix} q^4 \\ V \\ p = (q-1)n+1 \end{pmatrix}
$$
 for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

then

$$
[y_{iq}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

$$
([a_{ij}]\ \overline{\wedge} [c_{ik}]) = [y_{iq}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

$$
([x_{ij}]\ \underline{\vee} [y_{ik}]) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \underline{\vee} \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

_{SO}

 $[z_{ip}] \ = \ \left[\begin{array}{rrrrrrrr} 1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 1 \end{array} \right]$

where z_{ip} = max(x_{ij} , 1 - y_{ik}) such that $p = n(j - 1) + k$

if $h_{iq} = \left(\bigwedge_{p=(q-1)4+1}^{q4} (z_{ip}) \right)$ for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

$$
[h_{iq}] = ([x_{ij}] \vee [y_{ik}]) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

$$
([a_{ij}] \wedge [b_{ij}]) \vee ([a_{ij}] \wedge [c_{ij}]) = ([x_{ij}] \vee [y_{ik}]) = [h_{iq}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

$$
\text{so}[a_{ij}] \wedge ([b_{ij}] \vee [c_{ij}]) \neq ([a_{ij}] \wedge [b_{ij}]) \vee ([a_{ij}] \wedge [c_{ij}])
$$

3.1.25 \bf{Remark}

Š

Or-Not-B-Product is not distributive over And-Not-B-Product. $[a_{ij}] \vee (b_{ij}] \wedge [c_{ij}]) \neq ([a_{ij}] \vee [b_{ij}]) \wedge ([a_{ij}] \vee [c_{ij}])$

$3.1.26$ Remark

Let $[a_{ij}], [b_{ij}], \in SM_{m \times n}$ and $* \in (\wedge, \vee, \overline{\wedge}, \vee)$ be the binary operation. Then $[a_{ij}]$ * $[b_{ij}] \neq [b_{ij}] \ast [a_{ij}].$

3.1.27 Example

Let $U = \{u_1, u_2, u_3\}$ be the universal set, $E = \{e_1, e_2, e_3, e_4\}$ be the set of parameters and $A = \{e_1, e_2\}, B = \{e_3, e_4\}$ be the subsets of E.

Let $f_A: E \to P(U)$ be such that $f_A(e_1) = \{u_1, u_2\}$ $f_A(e_2) = \{u_2, u_3\}$ $f_A(e_3) = f_A(e_4) = \phi$ $R_A = \{(u_1, e_1), (u_2, e_2), (u_2, e_2), (u_3, e_2)\}\$

and $f_B: E \to P(U)$ be such that $f_B(e_3)=U$ $f_B(e_4) = \{u_1, u_3\}$ $f_B(e_1) = f_B(e_2) \doteq \phi$ $R_B = \{(u_1, e_3), (u_2, e_3), (u_3, e_3), (u_1, e_4), (u_3, e_4)\}\$

$$
B = [b_{ij}] = \left[\begin{array}{rrr} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{array} \right]
$$

$$
[\dot{a}_{ij}] \wedge [b_{ik}] = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \wedge \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

\n
$$
[d_{ip}] = \begin{bmatrix} 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}
$$

where $d_{ip} = \min(a_{ij}, b_{ik})$ such that $p = n(j-1) + k$

if
$$
y_{iq} = \begin{pmatrix} q4 \\ V \\ p=(q-1)4+1 \end{pmatrix}
$$
 for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

 ${\bf Then}$

Ý.

 $[Y] = [y_{iq}]_{3 \times 4} = [a_{ij}] \wedge [b_{ik}]$

$$
[a_{ij}] \wedge [b_{ik}] = \left[\begin{array}{rrr} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array}\right]
$$

Now

€

$$
[b_{ij}] \wedge [a_{ik}] = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix} \wedge \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}
$$

$$
[f_{ip}] = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \end{bmatrix}
$$

where $f_{ip} = \min(a_{ij}, b_{ik})$ such that $p = n(j-1) + k$

if
$$
e_{iq} = \begin{pmatrix} q4 \\ V \\ p=(q-1)4+1 \end{pmatrix}
$$
 for all $i = 1, 2, 3$ and $q = 1, 2, 3, 4$

Then

$$
[E] = [e_{iq}]_{3 \times 4} = [b_{ij}] \wedge [a_{ik}]
$$

$$
[b_{ij}] \wedge [a_{ik}] = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}
$$

hence

 $[a_{ik}] \wedge [b_{ij}] \neq [b_{ij}] \wedge [a_{ik}]$.

Therefore commutative law does not hold with respect to And-B-Product similarly

 $[a_{ik}] \vee [b_{ij}] \neq [b_{ij}] \vee [a_{ik}]$ $[a_{ik}] \barwedge [b_{ij}] \neq [b_{ij}] \barwedge [a_{ik}]$ $[a_{ik}] \veeeq [b_{ij}] \neq [b_{ij}] \veeeq [a_{ik}]$

3.1.28 Theorem

والمراجح

Let $SM_{m \times n}$ be the collection of all the soft matrices and $* \in \{\wedge, \vee, \overline{\wedge}, \vee\}$ be the binary operations, then $(SM_{m \times n}, *)$ is a semigroup.

Proof. Straightforward. ■

3.L.29 Theorem

 \mathcal{G}

Let $SM_{m \times n}$ be the collection of all the soft matrices and \ast , $o \in \{\wedge, \vee\}$ be the binary operations, Then $(SM_{m\times n}, *$, o) is a semiring.

Proof. Straightforward. \blacksquare

3.2 Soft Matrix Decision Making

In this section we construct a soft matrix decision making, with the help of soft matrix decision function and then select an optimum solution from the decision soft matrix.

3.2.1 Definition

let $[a_{ij}]$, $[b_{ij}] \in SM_{m \times n}$, and let $[c_{ij}]$ be the product of $[a_{ij}]$ and $[b_{ij}]$. Then the soft matrix decision function, denoted SMDF is define as follows

$$
SMDF:SM_{m\times n}\to SM_{m\times 1}
$$

$$
SMDF[c_{ij}] = \left[\frac{\sum_{j=1}^{n} \{c_{ij}\}}{n}\right] \text{ where } i = 1, 2, ..., m
$$

the one column soft matrix $SMDF[c_{ij}]$ is called decision soft matrix.

3.2.2 Definition

let $U = {u_1, u_2,...u_n}$ be initial universe and $SMDF[c_{ij}] = [d_{i1}]$. Then a subset of U can be obtained by using $[d_{i1}]$ as in he following way

 $\mathit{optm}_{[d_{i1}]}(U) = \{u_i : u_i \in U, \max(d_{i1})\}$

3.2.3 Applications

Assume that a person wants to seek admission in Ph.D. program and the universal set contain different universities $U = \{u_1, u_2, u_3, u_4, u_5\}$, which may be characterized by a set of parameters $E = \{e_1, e_2, e_3, e_4\}$. For $j = 1, 2, 3, 4$ the parameters e_j stand for "Part time studies", "less Fee ", "Full time studies" and "Located near Islamabad" respectively. Then we can give the following examples.

3.2.4 Example

v

 \tilde{f}

 \blacksquare^p

Suppose that two Students, Mr. A and Mr. B, come to the contact with each other and want to get admission. If each of them has to consider their own set of parameters, then we select a University on the basis of the sets of partners' parameters by using the Soft Matrix Decision as follows.

Assume that $U = \{u_1, u_2, u_3, u_4, u_5\}$ is a universal set and $E = \{e_1, e_2, e_3, e_4\}$ is a set of all parameters.

Mr. A and Mr. B have to choose the sets of their parameters, $A = \{e_2, e_3, e_4\}$ and $B = \{e_1, e_3, e_4\}$ respectively.

Then we can write the following soft matrices which are constructed according to their parameters.

Now, we can find a product of the soft matrices $[a_{ij}]$ and $[b_{ik}]$ by using And-Bproduct as follows

Now we apply And-B-product since both Mr. A and Mr. B choices have to be considered.

> $[d_{ip}]$ = 000000010 0000100110 0000001100110000 0000001100 0000000000

and $d_{ip} = \min(a_{ij}, b_{ik})$ such that $p = n(j - 1) + k$

Where $y_{iq} = \begin{pmatrix} \bigvee & (d_{ip}) \end{pmatrix}$ $\sqrt{p=(q-1)4+1}$ /

for all $i = 1, 2, 3, 4, 5$ and $q = 1, 2, 3, 4$

Then

 $[Y] = [y_{iq}]_{5 \times 4} = [a_{ij}] \wedge [b_{ik}]$

$$
[a_{ij}] \wedge [b_{ik}] = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{bmatrix}
$$

We can find a decision soft matrix as $\overline{5}$

$$
MDF([a_{ij}] \wedge [b_{ik}]) = \begin{bmatrix} 0.5 \\ 0.75 \\ 0.5 \\ 0.25 \\ 0.5 \end{bmatrix}
$$

 $\tilde{\mathbf{y}}$

we can find an optimum set of U according to $MDF([a_{ij}]\wedge [b_{ik}])^{i}$ $\varphi(t) = \{u_1\}$, where u_2 is an optimum University for Mr. A and Mr. B.

Note that the optimal set of U may contain more than one element.

Similarly, we can also use the other products $([a_{ij}] \vee [b_{ik}])$, $([a_{ij}] \wedge [b_{ik}])$ and $([a_{ij}] \vee [b_{ik}])$ for the other convenient problems.

Conclusion

Conclusion

The soft set theory has been used in different fields. The results of this thesis show that the B-products are binary. Further it is shown that associative laws as well as distributive laws holds. At the end of this thesis we highlighted that soft matrix decision making on the basis of soft set theory is useful. The example of a student who is looking for some university for Ph.D. is also given in this thesis. These type of products can also be defined in fizzy soft matrices. and we can.also take the products of the soft sets and then couvert it into soft matrices and can compair the result in both the cases. This Converse can be applied in both soft matrices and fuzzy soft matrices.

it

I .
I

Bibliography

 $\overline{\mathbb{R}}$

- [1] H. Aktas and N. Qa\$man, Soft sets and soft groups, Information Sciences ¹⁷⁷ (2007) 2726-2735
- [2] M. I. Ali, F. Feng, X. Liu, W.K. Min and M. Shabir, On some new operations in soft set theory, Computers and Mathematics with Applications 57 (2009) 1547-1553
- [3] M.I. Ali, M. Shabir and M. Naz, Algebraic structures of soft sets associated with new operations, Computers and Mathematics with Applications, 61 (9) (2011) 2647-2654.
- [4] M.I. Ali and M. Shabir, Logic connectives for soft sets and fuzzy soft sets, IEEE Transactions on 22 (6) (2014), 1431-1442.
- [5] N. Cagman and S. Enginoglu, Soft matrix theory and its decision making, Computers and Mathematics with Applications 59 (2010) 3308-3314.
- [6] D.chen, E. c. c. Tsang and D.s. Yeung, some notes on the parameterization reduction of soft sets, in: International Conference on Machine Learning and Cybernetics, vol. 3, (2003) 1442-1445.
- [7] D. Chen, E. C. C. Tsang, D. S. Yeung and X. Wang, The parameterization reduction of soft sets and its applications, Computers and Mathematics with Applications 49 (2005) 757-763.
- [8] F. Feng, Y. B. Jun and X. Zhao, Soft semirings, Computers and. Mathematics with Applications 56 (10) (2008) 2621-2628.
- [9] Y. B. Jun, Soft BCK/BCI-algebras, Computers and Mathematics with Applications 56 (2008) 1408-1413.
- [10] Y. B. Jun and C.H. Park, Applications of soft sets in ideal theory of BCK/BCIalgebras, Information Sciences 178 (2008) 2466-24Tb.

 $\overline{}$

- [25] A. R. Roy, P.K. Maji, A fuzzy soft set theoretic approach to decision making problems, Journal of computational and Applied Mathematics 203 (2007) 412-418.
- [26] Q. -M. Sun, Z. -L. Zhang and J. Liu, Soft sets and soft modules, in: Guoyin Wang Tian-rui Li, Jerzy W. Grzymala-Busse, Duoqian Miao, Andrzej Skowron, Yiyu Yao (Eds.), Rough Sets and Knowledge Technology, RSKT-2008, Proceedings, Springer, 2008 403-409.
- [27] Z. Xiao, Y. Li, B. Zhong and X. Yang, Research on synthetically evaluating method for business competitive capacity based on soft set, Statistical Research (2003) 52-54.
- [28] Z. Xiao, L. Chen, B. Zhong and S. Ye, Recognition for soft information based on the theory of soft sets, in: J. Chen (Ed.), Proceedings of ICSSSM-O5, vol. 2,IEEE, (2005) 1104-1106.
- [29] Z. Xiao, K. Gong and Y. Zou, A combined forecasting approach based on fuzzy soft sets, Journal of Computational and Applied Mathematics 228 (2009) 326-333.872-880.
- [30] X. Yang, D. Yu, J. Yang and C. Wu, Generalization of soft set theory: from crisp to fuzzy case, in: Bing-Yuan Cao (Ed.), Fuzzy Information and Engineering: Proceedings of ICFIE-2007, in: Advances in Soft Computing, vol. 40, Springer,(2007), 345-355. \sim
	- [31] Y. Zou and Z. Xiao, Data analysis approaches of soft sets under incomplete information, Knowledge-Based Systems 21 (2008) 941-945.

D

 $\widetilde{\mathcal{R}}$

\

Conclusion

The soft set theory has been used in different fields. The results of this thesis show that the B-products are binary. Further it is shown that associative laws as well as distributive laws holds. At the end of this thesis we highlighted that soft matrix decision making on the basis of soft set theory is useful. The example of a student who is looking for some university for Ph.D. is also given in this thesis. These type of products can also be defined in fuzzy soft matrices. and we can also take the products of the soft sets and then convert it into soft matrices and can compair the result in both the cases. This Converse cau be applied in both soft matrices and fuzzy soft matrices.

Bibliography

--F

- [1] H. Aktas and N. Ça \S man, Soft sets and soft groups, Information Sciences 177 (2007) 2726-2735
- [2] M. I. Ali, F. Feng, X. Liu, W.K. Min and M. Shabir, On some new operations in soft set theory, computers and Mathematics with Applications 57 (2009) 1547- 1553
- $[3]$ M.I. Ali, M. Shabir and M. Naz, Algebraic structures of soft sets associated wi new operations, computers and Mathematics with Applications, 61 (9) (2011) 2647-2654.
- [4] M.I. Ali and M. Shabir, Logic connectives for soft sets and fuzzy soft sets, IEEE Transactions on 22 (6) (2014), 1431-1442.
- [5] N. Cagman and S. Enginoglu, Soft matrix theory and its decision making, Computers and Mathematics with Applications 59 (2010) 3308-3314.
- [6] D.Chen, E. C. C. Tsang and D.S. Yeung, Some notes on the parameterization reduction of soft sets, in: International Conference on Machine Learning aud Cybernetics, vol. 3, (2003) 1442-1445.
- [7] D. Chen, E. C. C. Tsang, D. S. Yeung and X. Wang, The parameterization reduction of soft sets and its applications, Computers and Mathematics with Applications 49 (2005) 757-763'
- [8] F. Feng, Y. B. Jun and X. Zhao, Soft semirings, Computers and Mathematics with Applications 56 (10) (2008) 2621-2628.
- [g] y. B. Jun, Soft BCK/BCI-algebras, computers and Mathematics with Applications 56 (2008) 1408-1413.
- [10] Y. B. Jun and C.H. Park, Applications of soft sets in ideal theory of BCK/BCIalgebras, Information Sciences 178 (2008) 2466-2475.

j

kjI

IIlIrI

I ,II

'BLYJ

fri
1

ttr'l ll

[, 1 $k_{\tau_{\tau_{\tau}}}$ f r 1- lrt \mathbf{l} .

,
|
|

,T, : 1"1. t.l Irl l:, ril

\$ بر
مج '{-i l, iil ζ : rt iI i!' lir'+{ r\$'

ff iirli f,

ir d. s

t, E '-

- [25] A. R. Roy, P.K. Maji, A fuzzy soft set theoretic approach to decision making problems, Journal of computational and Applied Mathematics 203 (2007) 4L2- 418.
- [26] Q. -M. Sun, Z. -L. Zharg and J. Liu, Soft sets and soft modules, in: Guoyin Wang, Tian-rui Li, Jerzy W. Grzymala-Busse, Duoqian Miao, Andrzej Skowron, Yiyu Yao (Eds.), Rough Sets and Knowledge Technology, RSKT-2008, Proceedings, Springer, 2008 403-409.
- l27l Z. Xiao, Y. Li, B. Zhong and X. Yang, Research on synthetically evaluating method for business competitive capacity based on soft set, Statistical Research (2003) 52-54.
- 128) Z. Xiao, L. Chen, B. Zhong and S. Ye, Recognition for soft information based on the theory of soft sets, in: J. Chen (Ed.), Proceedings of ICSSSM-05, vol. 2, IEEE, (2005) 1104-1106. í.
- 129) Z. Xiao, K. Gong and Y. Zou, A combined forecasting approach based on fuzzy soft sets, Journal of Computational and Applied Mathematics 228 (2009) 326- 333.872-880.
- [30] X. Yang, D. Yu, J. Yang and C. Wu, Generalization of soft set theory: from crisp to fuzzy case, in: Bing-Yuan Cao (Ed.), Fuzzy Information and Engineering: Proceedings of ICFIB2007, in: Advances in Soft Computing, vol.40, Springer,(2007), 345-355. 418.

(26) Q. -M, Sun, Z. -L. Zhang and J. Liu, Soft sets and soft mo

Tian-rui Li, Jerzy W. Grzymala-Busse, Duoqian Miao,

Yao (Eds.), Rough Sets and Knowledge Technology, R.

Springer, 2008 403-409.

(27) Z. Xiao, Y. Li,
	- [31] Y. Zou and Z. Xiao, Data analysis approaches of soft sets under incomplete in-

 $\mathbf{\overline{1}}$ I